

ON NEIGHBOURLY IRREGULAR GRAPHS

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ABSTRACT. A connected graph G is said to be neighbourly irregular graph if no two adjacent vertices of G have same degree. In this paper we obtain neighbourly irregular subdivision graphs, line graphs and total graphs. The neighbourly irregularity of some graph products are also investigated.

1. INTRODUCTION

Throughout this paper we consider finite, simple, connected graphs. For graph theoretical terminology we follow the book [8]. Let G be a graph with n vertices and m edges. The vertex set and edge set of G are denoted by $V(G)$ and $E(G)$ respectively.

Let $\deg_G(v)$ denote the degree of a vertex v in G . A graph G is said to be regular if all its vertices have the same degree. A connected graph G is said to be highly irregular if each neighbor of any vertex has different degree [1]. It is called k -neighbourhood regular if each vertex is adjacent to exactly k vertices of the same degree [5]. The graph G is said to be neighbourly irregular graph, abbreviated as NI graph, if no two adjacent vertices of G have the same degree. This concept was introduced by Bhramsm and Ayyaswamy [6]. They constructed NI graphs of order n for a given n and a partition of n with distinct parts and proved some properties of NI graphs related to graphoidal covering number, gracefulness, ply number, lace number, clique graph and minimal edge covering. The Figure 1 depicts an example of NI graph.

We now present some graph operations that will be used in this paper. Let G and H be two graphs. The join $G + H$ of graphs G and H with disjoint vertex sets V_1 and V_2 and edge sets E_1 and E_2 is the graph union $G \cup H$ together with all the edges joining V_1 and V_2 . The Cartesian product of two graphs G and H is the graph $G \times H$

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whose vertex set is $V(G \times H) = V_1 \times V_2$ and two vertices (u_1, u_2) and (v_1, v_2) are adjacent in $G \times H$ whenever $u_1 = v_1$ and u_2 is adjacent to v_2 in H or $u_2 = v_2$ and u_1 is adjacent to v_1 in G . The corona product $G \circ H$ is obtained by taking one copy of G and $|V(G)|$ copies of H ; and by joining each vertex of the i -th copy of H to the i -th vertex of G , $i = 1, 2, \dots, |V(G)|$. The lexicographic product $G[H]$ of graphs G and H is the graph with vertex set $V(G) \times V(H)$ and $u = (u_1, v_1)$ is adjacent with $v = (u_2, v_2)$ whenever $(u_1$ is adjacent to $u_2)$ or $(u_1 = u_2$ and v_1 is adjacent to $v_2)$. The tensor product of the graphs G and H is the graph $G \otimes H$ whose vertex set is $V(G \otimes H) = V_1 \times V_2$ and two vertices (u_1, u_2) and (v_1, v_2) are adjacent in $G \otimes H$ if and only if u_1 is adjacent to v_1 in G and u_2 is adjacent to v_2 in H , see for details [7].

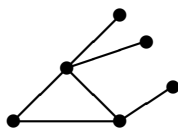


Figure 1: NI Graph.

Let G and H be graphs. The strong product $G \boxtimes H$ of graphs G and H is the graph with vertex set $V(G) \times V(H)$ and $u = (u_1, v_1)$ is adjacent with $v = (u_2, v_2)$ whenever $(v_1 = v_2$ and u_1 is adjacent with $u_2)$ or $(u_1 = u_2$ and v_1 is adjacent with $v_2)$ or $(u_1$ is adjacent with u_2 and v_1 is adjacent with $v_2)$, see [11] for details.

A graph G with a specified vertex subset $U \subseteq V(G)$ is denoted by $G(U)$. Suppose G and H are graphs and $U \subseteq V(G)$. The generalized hierarchical product, denoted by $G(U) \square H$, is the graph with vertex set $V(G) \times V(H)$ and two vertices (g, h) and (g', h') are adjacent if and only if $g = g' \in U$ and $hh' \in E(H)$ or, $gg' \in E(G)$ and $h = h'$. This graph operation introduced recently by Barri ere et al. [3, 4] and found some applications in computer science. We encourage the reader to consult [2, 10] for mathematical properties and chemical applications of the hierarchical product of graphs.

In this paper we obtain neighbourly irregular subdivision graphs, line graphs and total graphs. The neighbourly irregular graphs obtained from some graph operations are also considered into account.

2. NEIGHBOURLY IRREGULAR GRAPHS

The subdivision graph $S(G)$ of a graph G is obtained from G by inserting a new vertex onto every edge of G [8].

Theorem 2.1. *Let G be a graph. The subdivision graph $S(G)$ is NI if and only if G does not have any vertex of degree two.*

Proof. Suppose G does not have any vertex of degree two, i.e. for every vertex $u \in V(G)$, $\deg_G(u) \neq 2$. If $u \in V(G)$ then $\deg_G(u) = \deg_{S(G)}(u)$. In $S(G)$ for every edge $e = uv$ in G , a new vertex w is inserted and $\deg_{S(G)}(w) = 2$. Therefore

$\deg_{S(G)}(u) \neq \deg_{S(G)}(w)$ for every adjacent pair (u, w) of the vertices of $S(G)$. Hence $S(G)$ is NI graph. Conversely, if G is a graph with at least one vertex of degree two then $S(G)$ is not NI graph. \square

The line graph $L(G)$ of a graph G is the graph whose vertices corresponds to the edges of G and two vertices in $L(G)$ are adjacent if and only if the corresponding edges are adjacent in G [8]. If $e = uv$ is an edge of G then $\deg_{L(G)}(e) = \deg_G(u) + \deg_G(v) - 2$. Let $N(u)$ denotes the set of vertices which are adjacent to u .

Theorem 2.2. *For any graph G , its line graph $L(G)$ is NI graph if and only if $N(u)$ contains all vertices of different degree for all $u \in V(G)$.*

Proof. Let $L(G)$ be NI graph. To prove that $N(u)$ contains all vertices of different degree for all $u \in V(G)$, on the contrary, suppose $N(u)$ contains two vertices v and w of same degree, that is, $\deg_G(v) = \deg_G(w)$, where $v, w \in N(u)$. Therefore,

$$\deg_G(u) + \deg_G(v) - 2 = \deg_G(u) + \deg_G(w) - 2.$$

This implies that, $\deg_{L(G)}(e_1) = \deg_{L(G)}(e_2)$, where $e_1 = uv$ and $e_2 = uw$. Thus, $L(G)$ is not NI graph, a contradiction. Hence $N(u)$ contains all vertices of different degree.

Conversely, let $\deg_G(v) \neq \deg_G(w)$ for all $v, w \in N(u)$ and $u \in V(G)$. Therefore,

$$\deg_G(u) + \deg_G(v) - 2 \neq \deg_G(u) + \deg_G(w) - 2.$$

That is, $\deg_{L(G)}(e_1) \neq \deg_{L(G)}(e_2)$, where $e_1 = uv$ and $e_2 = uw$. Hence $L(G)$ is NI graph. \square

Lemma 2.1. ([6]) *If u is a vertex of maximum degree in a NI graph then $N(u)$ contains at least two vertices of same degree.*

Theorem 2.3. *If G is NI graph then $L(G)$ is not NI graph.*

Proof. Let u be the vertex of maximum degree in a NI graph G . Then by Lemma 2.1, $N(u)$ contains at least two vertices say v and w of same degree. The vertices $e_1 = uv$ and $e_2 = uw$ are adjacent in $L(G)$. But,

$$\begin{aligned} \deg_{L(G)}(e_1) &= \deg_G(u) + \deg_G(v) - 2 \\ &= \deg_G(u) + \deg_G(w) - 2 \\ &= \deg_{L(G)}(e_2). \end{aligned}$$

Hence $L(G)$ is not NI graph. \square

Theorem 2.4. *For each integer $k \geq 1$, there exists a graph G with maximum degree $\Delta(G) = k$ such that $L(G)$ is NI graph.*

Proof. Let u and v be two vertices of G such that, $\deg_G(u) = \deg_G(v) = k$. Let the vertices u_1, u_2, \dots, u_{k-1} and v be adjacent to u and the vertices v_1, v_2, \dots, v_{k-1} and u be adjacent to v , see Figure 2.

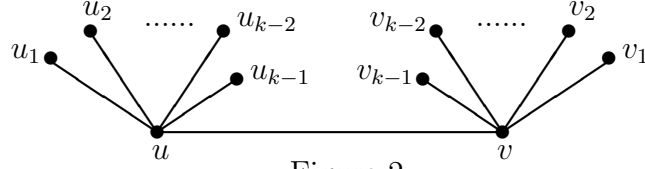


Figure 2

Join the vertices u_{k-i} to the vertices $v_{i+1}, v_{i+2}, \dots, v_{k-1}$ for $i = 1, 2, \dots, k-2$. Thus the graph G is obtained. It is not difficult to see that $\deg_G(u_i) = i$ for $i = 1, 2, \dots, k-1$ and $\deg_G(v_i) = i$ for $i = 1, 2, \dots, k-1$ and $\deg_G(u) = \deg_G(v) = k$.

Now $N(u_i) = \{u, v_{k-i+1}, v_{k-i+2}, \dots, v_{k-1}\}$ and $\deg_G(u) = k$ and $\deg_G(v_{k-i+1}) = k - i + 1$ for $i = 2, 3, \dots, k-1$. Hence $N(u_i)$ contains all vertices of different degrees, $i = 1, 2, \dots, k-1$. Similarly we can show that $N(v_i)$ contains all vertices of different degrees for $i = 1, 2, \dots, k-1$. Also $N(u)$ and $N(v)$ contains all vertices of different degrees. Hence by Theorem 2.2, $L(G)$ is NI graph. \square

The total graph $T(G)$ of a graph G is a graph whose vertex set is $V(G) \cup E(G)$ and two vertices in $T(G)$ are adjacent if and only if the corresponding elements are adjacent or incident in G [8]. If u is a vertex of G , then $\deg_{T(G)}(u) = 2 \deg_G(u)$. If $e = uv$ is an edge of G then $\deg_{T(G)}(e) = \deg_G(u) + \deg_G(v)$.

Theorem 2.5. *For any graph G , the total graph $T(G)$ is NI graph if and only if both G and $L(G)$ are NI graphs.*

Proof. Let the vertex set be $V(G) = \{v_1, v_2, \dots, v_n\}$ and the edge set be $E(G) = \{e_1, e_2, \dots, e_m\}$. Suppose G and $L(G)$ are NI graphs. In $T(G)$, let $e = xy$ be an edge then $x, y \in V(G)$ or $x, y \in E(G)$ or $x \in V(G)$ and $y \in E(G)$.

- (a) $x, y \in V(G)$. Since $\deg_G(x) \neq \deg_G(y)$, $\deg_{T(G)}(x) = 2 \deg_G(x) \neq 2 \deg_G(y) = \deg_{T(G)}(y)$.
- (b) $x, y \in E(G)$. Let $x = v_i v_j$ and $y = v_i v_k$, so that x and y are adjacent in $T(G)$. Therefore $\deg_{T(G)}(x) = \deg_G(v_i) + \deg_G(v_j)$ and $\deg_{T(G)}(y) = \deg_G(v_i) + \deg_G(v_k)$. But $\deg_{L(G)}(x) \neq \deg_{L(G)}(y)$ as $L(G)$ is NI graph. And $\deg_{L(G)}(x) = \deg_G(v_i) + \deg_G(v_j) - 2$ and $\deg_{L(G)}(y) = \deg_G(v_i) + \deg_G(v_k) - 2$. Therefore $\deg_{T(G)}(x) \neq \deg_{T(G)}(y)$.
- (c) $x \in V(G)$ and $y \in E(G)$. Let $e = xy = v_i e_j$ for some $v_i \in V(G)$ and $e_j \in E(G)$. Therefore $\deg_{T(G)}(x) = \deg_{T(G)}(v_i) = 2 \deg_G(v_i)$ and

$$\begin{aligned}
 \deg_{T(G)}(y) &= \deg_{T(G)}(e_j) = \deg_{L(G)}(e_j) + 2 \\
 &= \deg_G(v_i) + \deg_G(v_k) - 2 + 2, \text{ where } e_j = v_i v_k \\
 &= \deg_G(v_i) + \deg_G(v_k) \\
 &\neq \deg_G(v_i) + \deg_G(v_i) \text{ since } G \text{ is NI graph, } \deg_G(v_i) \neq \deg_G(v_k) \\
 &= 2 \deg_G(v_i) \\
 &= \deg_{T(G)}(x).
 \end{aligned}$$

Thus in all cases $T(G)$ is NI graph.

Conversely, suppose $T(G)$ is NI graph. We have to prove that both G and $L(G)$ are NI graphs. If G is not NI graph, then there exists an edge $e_k = v_i v_j$ in G such that $\deg_G(v_i) = \deg_G(v_j)$. So, $\deg_{T(G)}(v_i) = \deg_{T(G)}(v_j)$. A contradiction to $T(G)$ is NI graph.

Suppose $L(G)$ is not NI graph, then there exists two adjacent vertices $e_i = v_r v_s$ and $e_j = v_r v_k$ in $L(G)$ with $\deg_{L(G)}(e_i) = \deg_{L(G)}(e_j)$. Thus, $\deg_G(v_r) + \deg_G(v_s) - 2 = \deg_G(v_r) + \deg_G(v_k) - 2$. Hence $\deg_G(v_r) + \deg_G(v_s) = \deg_G(v_r) + \deg_G(v_k)$ and so $\deg_{T(G)}(e_i) = \deg_{T(G)}(e_j)$. Again a contradiction to $T(G)$ is NI graph. \square

From Theorems 2.3 and 2.5 we have following result:

Theorem 2.6. *There is no nontrivial graph G whose total graph $T(G)$ is NI graph.*

3. NEIGHBOURLY IRREGULAR GRAPH PRODUCTS

Let G_1 be the graph with vertex set $V_1 = V(G_1)$ and edge set $E_1 = E(G_1)$. Let G_2 be the another graph with vertex set $V_2 = V(G_2)$ and edge set $E_2 = E(G_2)$. If $(u, v) \in V(G_1 \times G_2)$ then $\deg_{G_1 \times G_2}((u, v)) = \deg_{G_1}(u) + \deg_{G_2}(v)$.

Theorem 3.1. *For any two graphs G and H , the Cartesian product $G \times H$ is NI graph if and only if both G and H are NI graphs.*

Proof. Let G and H be the NI graphs. Let $N(u) = \{u_1, u_2, \dots, u_k\}$ and $N(v) = \{v_1, v_2, \dots, v_l\}$, where $u \in V(G)$ and $v \in V(H)$. Then by definition of NI graph $\deg_G(u) \neq \deg_G(u_i)$ for $i = 1, 2, \dots, k$ and $\deg_H(v) \neq \deg_H(v_j)$ for $j = 1, 2, \dots, l$. The vertex set of $G \times H$ is $V(G) \times V(H)$. Clearly,

$$N((u, v)) = \{(u_i, v), (u, v_j) | i = 1, 2, \dots, k \text{ and } j = 1, 2, \dots, l\}.$$

To prove that $G \times H$ is NI graph, on the contrary $G \times H$ is not NI graph. Then there exists at least one vertex in $N((u, v))$ with same degree as the degree of the vertex (u, v) in $G \times H$. Hence $\deg_{G \times H}((u, v)) = \deg_{G \times H}((u_i, v))$ and so $\deg_G(u) + \deg_H(v) = \deg_G(u_i) + \deg_H(v)$. This implies that $\deg_G(u) = \deg_G(u_i)$. A contradiction to the fact that G is NI graph. Similarly we can show that a contradiction to the fact that H is NI graph. Hence $G \times H$ is NI graph.

Conversely, suppose $G \times H$ is NI graph. Let $N(u) = \{u_1, u_2, \dots, u_k\}$ and $N(v) = \{v_1, v_2, \dots, v_l\}$, where $u \in V(G)$ and $v \in V(H)$. Therefore,

$$\deg_{G \times H}((u, v)) \neq \deg_{G \times H}((u_i, v)),$$

where $(u_i, v) \in N((u, v))$. This shows that $\deg_G(u) + \deg_H(v) \neq \deg_G(u_i) + \deg_H(v)$ and so $\deg_G(u) \neq \deg_G(u_i)$. Hence G is NI graph. Similarly we can show that H is NI graph. \square

In the next theorem, the tensor product of graphs are considered. We notice that if $(u, v) \in V(G \otimes H)$ then $\deg_{G \otimes H}((u, v)) = \deg_G(u) \deg_H(v)$.

Theorem 3.2. *For any NI graph G and a regular graph H , the tensor product $G \otimes H$ is NI graph.*

Proof. Let G be the NI graph and H be any regular graph of degree r . Let $N(u) = \{u_1, u_2, \dots, u_k\}$ where $u \in V(G)$ and $N(v) = \{v_1, v_2, \dots, v_r\}$, where $v \in V(H)$. Since G is NI graph, $\deg_G(u) \neq \deg_G(u_i)$, $i = 1, 2, \dots, k$.

On the contrary, suppose $G \otimes H$ is not NI graph, then there exists at least two adjacent vertices of same degree in $G \otimes H$. Without loss of generality, suppose $\deg_{G \otimes H}((u, v)) = \deg_{G \otimes H}((u_i, v_j))$, where $i = 1, 2, \dots, k$ and $j = 1, 2, \dots, r$. Hence, $\deg_G(u) \deg_H(v) = \deg_G(u_i) \deg_H(v_j)$ and so $\deg_G(u) \deg_H(v) = \deg_G(u_i) \deg_H(v)$. Since H is a regular graph, $\deg_H(v) = \deg_H(v_j) = r$ and so $\deg_G(u) = \deg_G(u_i)$, a contradiction to the fact that G is NI graph. Hence $G \otimes H$ is NI graph. \square

Suppose G is a graph. The set of all vertex degrees of G is called the vertex degree set of G .

Theorem 3.3. *Let G and H be graphs with vertex degree sets S_G and S_H , respectively. Then $G + H$ is NI graph if and only if both G and H are NI graphs and, ($|V(G)| = |V(H)|$ and $S_G \cap S_H = \emptyset$) or ($|V(G)| - |V(H)| = t > 0$ and $S_G \cap \{a_i + t \mid a_i \in S_H\} = \emptyset$).*

Proof. If both G and H are NI graphs and, $|V(G)| = |V(H)|$ and $S_G \cap S_H = \emptyset$, or, $|V(G)| - |V(H)| = t > 0$ and $S_G \cap \{a_i + t \mid a_i \in S_H\} = \emptyset$, it is clear that $G + H$ is NI graph. So, it is enough to prove that if $G + H$ is NI graph, then both G and H are NI graphs and, $|V(G)| = |V(H)|$ and $S_G \cap S_H = \emptyset$, or, $|V(G)| - |V(H)| = t > 0$ and $S_G \cap \{a_i + t \mid a_i \in S_H\} = \emptyset$. By definition of $G + H$, it is clear that if $G + H$ is NI graph, then both G and H are NI graphs. Now, there are two separate cases as follows:

- (a) $|V(G)| = |V(H)|$. Let $u \in V(G)$. So, by definition of join of two graphs, u is adjacent to each vertex in copy of H in $G + H$ and since $G + H$ is NI graph, therefore, $\deg_{G+H}(u) - \deg_{G+H}(v) = \deg_G(u) - \deg_H(v) \neq 0$, for each v in copy of H . This implies that $S_G \cap S_H = \emptyset$.
- (b) $|V(G)| - |V(H)| = t > 0$. Suppose $u \in V(G)$, so again by definition of join of two graphs, u is adjacent to each vertex in the copy of H in $G + H$ and since $G + H$ is NI graph, thus, $\deg_{G+H}(u) - \deg_{G+H}(v) = \deg_G(u) + |V(H)| - \deg_H(v) - |V(G)| = \deg_G(u) - \deg_H(v) - t \neq 0$, for each v in the copy of H . Therefore, $\deg_G(u) \neq \deg_H(v) + t$, for each v in copy of H . This means that $S_G \cap \{a_i + t \mid a_i \in S_H\} = \emptyset$.

By Cases (a) and (b), the proof is completed. \square

To prove the next result, we have to present some notations. Let G' be a copy of G and H_i be the i -th copy of H in $G \circ H$, $1 \leq i \leq |V(G)|$. Then, $G \circ H$ is obtained by joining each vertex of the i -th copy of H to the i -th vertex (x_i) of G . A vertex of $G \circ H$ corresponding to the vertex u in H is denoted by u' . Also, we denote a vertex of $G \circ H$ corresponding to the vertex v in G by v' .

Theorem 3.4. *Let G and H be nontrivial graphs. Then $G \circ H$ is NI graph if and only if both G and H are NI graphs.*

Proof. Let G and H be NI nontrivial graphs. Then it is clear that $G \circ H$ is NI graph. Conversely, let G and H be two nontrivial graphs and $G \circ H$ is NI graph. Suppose $u'v' \in E(G \circ H)$ such that $u', v' \in V(H_i)$, then $\deg_{G \circ H}(u') - \deg_{G \circ H}(v') = \deg_H(u) - \deg_H(v) \neq 0$ and so H is NI graph. On the other hand, if $u'v' \in E(G \circ H)$ such that $u', v' \in V(G')$, then $\deg_{G \circ H}(u') - \deg_{G \circ H}(v') = \deg_G(u) - \deg_G(v) \neq 0$. Therefore, G is NI graph. \square

Theorem 3.5. *Let G and H be nontrivial graphs. Then $G[H]$ is NI graph if and only if both G and H are NI graphs.*

Proof. By definition of lexicographic product $G[H]$ we have that $\deg_{G[H]}((u, v)) = \deg_G(u)|V(H)| + \deg_H(v)$ for each $(u, v) \in V(G[H])$. Thus, if G and H are NI graphs, it is clear that $G[H]$ is NI graph. Conversely, let $G[H]$ be NI graph. So, for every $(u, v_1)(u, v_2) \in E(G[H])$ that $v_1v_2 \in E(H)$ and $u \in V(G)$, $\deg_{G[H]}((u, v_1)) - \deg_{G[H]}((u, v_2)) = \deg_H(v_1) - \deg_H(v_2) \neq 0$, that is $\deg_H(v_1) \neq \deg_H(v_2)$ for every $v_1v_2 \in E(H)$. Therefore, H is NI graph. On the other hand, for every $(u_1, v)(u_2, v) \in E(G[H])$ that $u_1u_2 \in E(G)$ and $v \in V(H)$, $\deg_{G[H]}((u_1, v)) - \deg_{G[H]}((u_2, v)) = |V(H)|(\deg_G(u_1) - \deg_G(u_2)) \neq 0$. This means that $\deg_G(u_1) \neq \deg_G(u_2)$ for every $u_1u_2 \in E(G)$. Therefore, G is NI graph, which completes the proof. \square

Theorem 3.6. *Let G and H be nontrivial graphs and let U be a nonempty subset of $V(G)$. Then $G(U) \sqcap H$ is NI graph if and only if H is NI graph and for each $g_1g_2 \in E(G)$ and $v \in V(H)$,*

$$\deg_G(g_2) - \deg_G(g_1) \neq \begin{cases} 0, & \text{if } g_1, g_2 \in U \text{ or } g_1, g_2 \in V(G) - U, \\ \deg_H(v), & \text{if } g_1 \in U \text{ and } g_2 \in V(G) - U. \end{cases}$$

Proof. Let G and H be two nontrivial graphs and let U be a nonempty subset of $V(G)$. If $g \in U$ and $h \in V(H)$, then by definition of generalized hierarchical product, $\deg_{G(U) \sqcap H}((g, h)) = \deg_G(g) + \deg_H(h)$. Similarly, if $g \in V(G) - U$, then $\deg_{G(U) \sqcap H}((g, h)) = \deg_G(g)$. Thus, if H is NI graph and the inequality given in the statement of this theorem is satisfied, then clearly $G(U) \sqcap H$ is NI graph. Conversely, let $G(U) \sqcap H$ be NI graph. Suppose $(g_1, h_1)(g_2, h_2) \in E(G(U) \sqcap H)$. If $g_1 = g_2 \in U$ and $h_1h_2 \in E(H)$, then $\deg_{G(U) \sqcap H}((g_1, h_1)) - \deg_{G(U) \sqcap H}((g_1, h_2)) = \deg_H(h_1) - \deg_H(h_2)$ and since $G(U) \sqcap H$ is NI graph, thus $\deg_H(h_1) - \deg_H(h_2) \neq 0$. This implies that H is NI graph. On the other hand, if $g_1g_2 \in E(G)$ and $h_1 = h_2 \in V(H)$, we have two possible cases as follows:

- (a) $g_1 \in U$ and $g_2 \in V(G) - U$. So, $\deg_{G(U) \sqcap H}((g_1, h_1)) - \deg_{G(U) \sqcap H}((g_2, h_2)) = \deg_G(g_1) + \deg_H(h_1) - \deg_G(g_2)$ and since $G(U) \sqcap H$ is NI graph, $\deg_G(g_1) + \deg_H(h_1) - \deg_G(g_2) \neq 0$. This means that $\deg_G(g_2) - \deg_G(g_1) \neq \deg_H(h_1)$.
- (b) $g_1, g_2 \in U$ or $g_1, g_2 \in V(G) - U$. In this case, we have $\deg_{G(U) \sqcap H}((g_1, h_1)) - \deg_{G(U) \sqcap H}((g_2, h_2)) = \deg_G(g_1) - \deg_G(g_2)$ and since $G(U) \sqcap H$ is NI graph,

$\deg_G(g_1) - \deg_G(g_2) \neq 0$. This implies that $\deg_G(g_1) - \deg_G(g_2) \neq 0$, for each $g_1g_2 \in E(G)$ such that $g_1, g_2 \in U$ or $g_1, g_2 \in V(G) - U$.

This completes the proof. \square

Theorem 3.7. *If G and H are NI nontrivial graphs and there are no edges $g_1g_2 \in E(G)$, $h_1h_2 \in E(H)$ such that $\deg_G(g_1) = \deg_H(h_1)$ and $\deg_G(g_2) = \deg_H(h_2)$, then $G \boxtimes H$ is NI graph.*

Proof. Let G and H be NI nontrivial graphs and there be no edges $g_1g_2 \in E(G)$, $h_1h_2 \in E(H)$ such that $\deg_G(g_1) = \deg_H(h_1)$ and $\deg_G(g_2) = \deg_H(h_2)$. If $u \in V(G)$ and $v \in V(H)$, then by definition of strong product, $\deg_{G \boxtimes H}((u, v)) = \deg_H(v) + \deg_G(u) + \deg_G(u)\deg_H(v)$. Suppose $(g_1, h_1)(g_2, h_2) \in E(G \boxtimes H)$. If $g_1 = g_2$ or $h_1 = h_2$, it is clear that $\deg_{G \boxtimes H}((g_1, h_1)) - \deg_{G \boxtimes H}((g_2, h_2)) \neq 0$. On the other hand, if $g_1 \neq g_2$ and $h_1 \neq h_2$, so $g_1g_2 \in E(G)$ and $h_1h_2 \in E(H)$. Since G and H are NI graphs and $\deg_G(g_1) \neq \deg_H(h_2)$ and $\deg_G(g_2) \neq \deg_H(h_1)$, thus $\deg_{G \boxtimes H}((g_1, h_1)) - \deg_{G \boxtimes H}((g_2, h_2)) \neq 0$. This completes the proof. \square

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