

PSEUDO-BCK ALGEBRAS DERIVED FROM DIRECTOIDS

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ABSTRACT. The aim of this paper is to derive pseudo-BCK algebras from directoids and vice versa. We generalize some results proved by Ivan Chajda et al. in the case of BCK-algebras. We assign to an arbitrary pseudo-BCK algebra a semilattice-like structure and observe that this is the point where directoids are different from the semilattice-like structures. Finally, the relation between commutative deductive systems and derive directoids from a bounded pseudo-BCK(pDN) algebras and a characterization of commutative deductive systems of a bounded pseudo-BCK(pDN) algebra in terms of directoids is discussed.

1. INTRODUCTION

BCK-algebras were introduced by Y. Imai and K. Iséki in 1966 ([15, 19]) as algebras with a binary operation $*$ modeling the set-theoretical difference and with a constant element 0 that is a least element. S. Tanaka defined a special class of BCK-algebras called commutative BCK-algebras in 1975 (see [31]). In BCK-algebras, some lattices, as bounded commutative BCK-algebras, involutive BCK-lattices and bounded implicative BCK-algebras were defined and among the relationship between them were discussed [23]. Some recent researchers led to generalizations of the notion of pseudo structure on some types of algebras. G. Georgescu et al. [10] and independently J. Rachůnek [24], introduced pseudo-MV algebra which is a non-commutative generalization of MV-algebra. After a pseudo-MV algebra, the pseudo-BL algebra [11], the pseudo-BCK algebra [12] and as a generalization of BCI-algebra, the notion of pseudo-BCI algebra is introduced by W. A. Dudek et al. in [9]. A. Walendziak [32] introduced pseudo-BCH algebras as an extension of BCH-algebras. Further, he

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proved that every branchwise commutative pseudo-BCH algebra is a pseudo-BCI algebra [33]. Commutative pseudo-BCK algebras were originally defined by G. Georgescu et al. in [12] under the name of *semilattice-ordered pseudo-BCK algebras* and some properties of these structures were investigated by J. Kühr in [21, 22]. R. A. Borzooei et al. introduced in [1] (see also [2, 26, 27]) a pseudo-BE algebra as generalization of BE-algebra, and the commutative pseudo-BE algebra have recently been investigated by L. C. Ciungu. It was proved that the class of commutative pseudo-BE algebras is equivalent to the class of commutative pseudo-BCK algebras. Based on this result, all results holding for commutative pseudo-BCK algebras also hold for commutative pseudo-BE algebras [5]. Then she gave a characterization of commutative pseudo-BCK algebras and defined the commutative deductive systems of pseudo-BCK algebras and proved that a pseudo-BCK algebra \mathfrak{X} is commutative if and only if all the deductive systems of \mathfrak{X} are commutative. Also, she showed that the class of commutative pseudo-BCK algebras is a variety [6] (see also, [14]). A. Rezaei et al. introduced the notion of pseudo-CI algebras as an extension of pseudo-BE algebras and proved that the class of commutative pseudo-CI algebras coincide with the class of commutative pseudo-BCK algebras [28]. G. Georgescu et al. proved that every Wajsberg pseudo-hoop is a basic pseudo-hoop and every simple basic pseudo-hoop is a linearly ordered Wajsberg pseudo-hoop [13]. L. C. Ciungu in [7] showed that every pseudo-hoop is a pseudo-BCK-meet semilattice. The relation between FL_w -algebras, bounded pseudo-BCK(pP) algebras, pseudo-MTL algebras, pseudo-BL algebras and pseudo-MV algebras proved in [16]. Also, in [29, 30], the interrelationships between dual pseudo-Q/QC algebras and other pseudo algebras are visualized with a diagram and then they introduced the concepts of branchwise commutative pseudo-CI algebras and pointed pseudo-CI algebras and investigated some of properties. A. Iorgulescu for the first time introduced the notation of quasi-pseudo-M algebras as generalizations of pseudo-M algebras and (involutive) quasi-implicative-groups and the (strong involutive) (super) quasi-implicative-hoops, as generalizations of implicative-groups and implicative-hoops, respectively in [18]. I. Chajda et al. showed that one can be assign to an arbitrary BCK-algebra a semilattice-like structure every section of which possesses a certain antitone mappings [3], it arises a natural question of generalization of these concepts also for pseudo-BCK algebras. Since lattice theory has many applications in computer science and has an important and vital role in investigating the structure of a logical system, this motivated our investigations on directoids and pseudo-BCK-algebras to characterized several of its important properties. The main result of this paper establishes a bijective correspondence between pseudo-BCK algebras and some algebraic structures defined by two directoids. A characterization of commutative deductive systems of a bounded pseudo-BCK(pDN) algebra in terms of directoids is discussed and various results obtained mentioned in this paper can be transferred to the pseudo-BCK algebras. The core of the paper is based on by presenting a survey of some results of logic in the non-commutative case (see [3] for the commutative case) and extension of [25] (see also [4]).

2. PRELIMINARIES

In this section we recall some basic notions and results regarding (commutative) pseudo-BCK algebras.

Definition 2.1 ([9, 17]). An algebra $\mathfrak{X} = (X; \rightarrow, \rightsquigarrow, 1)$ of type $(2, 2, 0)$ is called a *pseudo-BCI algebra* if it satisfies the following axioms for all $x, y, z \in X$:

- (psBCI₁) $(x \rightarrow y) \rightsquigarrow ((y \rightarrow z) \rightsquigarrow (x \rightarrow z)) = 1$ and
 $(x \rightsquigarrow y) \rightarrow ((y \rightsquigarrow z) \rightarrow (x \rightsquigarrow z)) = 1$;
- (psBCI₂) $x \rightarrow ((x \rightarrow y) \rightsquigarrow y) = 1$ and $x \rightsquigarrow ((x \rightsquigarrow y) \rightarrow y) = 1$;
- (psBCI₃) $x \rightarrow x = x \rightsquigarrow x = 1$;
- (psBCI₄) $x \rightarrow y = y \rightsquigarrow x = 1 \Rightarrow x = y$;
- (psBCI₅) $x \preceq y$ if and only if $x \rightarrow y = 1$ if and only if $x \rightsquigarrow y = 1$.

A *pseudo-BCK algebra* [20] is a pseudo-BCI algebra $\mathfrak{X} = (X; \rightarrow, \rightsquigarrow, 1)$ satisfying the condition (psBCK), for all $x \in X$:

$$(psBCK) \quad x \rightarrow 1 = 1.$$

I. Chajda et al. proved that for every pseudo-BCI algebra $x \rightarrow y = 1$ if and only if $x \rightsquigarrow y = 1$ (see [4, Lemma 2.1]).

Remark 2.1. If $\mathfrak{X} = (X; \rightarrow, \rightsquigarrow, 1)$ is a pseudo-BCI algebra satisfying $x \rightarrow y = x \rightsquigarrow y$, for all $x, y \in X$, then $\mathfrak{X} = (X; \rightarrow, 1)$ is a BCI-algebra. Hence, every BCI-algebra is a pseudo-BCI algebra in a natural way.

Remark 2.2. By definition (psBCI₁)-(psBCI₅), pseudo-BCK algebras are contained in the class of pseudo-BCI algebras. A pseudo-BCI algebra which is not a pseudo-BCK algebra will be called *proper*.

From now on, \mathfrak{X} is a pseudo-BCK algebra, unless it is stated.

Proposition 2.1 ([12, 17]). *In any pseudo-BCK algebra \mathfrak{X} the following conditions hold for all $x, y, z \in X$:*

- (1) $x \preceq y$ implies $z \rightarrow x \preceq z \rightarrow y$ and $z \rightsquigarrow x \preceq z \rightsquigarrow y$;
- (2) $x \preceq y$ implies $y \rightarrow z \preceq x \rightarrow z$ and $y \rightsquigarrow z \preceq x \rightsquigarrow z$;
- (3) $x \rightarrow y \preceq (z \rightarrow x) \rightsquigarrow (z \rightarrow y)$ and $x \rightsquigarrow y \preceq (z \rightsquigarrow x) \rightarrow (z \rightsquigarrow y)$;
- (4) $x \rightarrow (y \rightsquigarrow z) = y \rightsquigarrow (x \rightarrow z)$ and $x \rightsquigarrow (y \rightarrow z) = y \rightarrow (x \rightsquigarrow z)$;
- (5) $x \preceq y \rightarrow x$ and $x \preceq y \rightsquigarrow x$;
- (6) $((x \rightarrow y) \rightsquigarrow y) \rightarrow y = x \rightarrow y$ and $((x \rightsquigarrow y) \rightarrow y) \rightsquigarrow y = x \rightsquigarrow y$.

Theorem 2.1 ([6]). *Let \mathfrak{X} be a pseudo-BCK algebra. The following statements are equivalent for all $x, y \in X$:*

- (1) \mathfrak{X} is commutative;
- (2) $x \rightarrow y = ((y \rightarrow x) \rightsquigarrow x) \rightarrow y$ and $x \rightsquigarrow y = ((y \rightsquigarrow x) \rightarrow x) \rightsquigarrow y$;
- (3) $(x \rightarrow y) \rightsquigarrow y = (((x \rightarrow y) \rightsquigarrow y) \rightarrow x) \rightsquigarrow x$ and
 $(x \rightsquigarrow y) \rightarrow y = (((x \rightsquigarrow y) \rightarrow y) \rightsquigarrow x) \rightarrow x$;
- (4) $x \preceq y$ implies $y = (y \rightarrow x) \rightsquigarrow x = (y \rightsquigarrow x) \rightarrow x$.

Definition 2.2 ([16]). If there is an element 0 of a pseudo-BCK algebra \mathfrak{X} , such that $0 \preceq x$ (i.e., $0 \rightarrow x = 0 \rightsquigarrow x = 1$), for all $x \in X$, then 0 is called the *zero* of \mathfrak{X} . A pseudo-BCK algebra with zero is called *bounded pseudo-BCK algebra* and it is denoted by $\mathfrak{X} = (X; \rightarrow, \rightsquigarrow, 0, 1)$.

Definition 2.3 ([16]). A *pseudo-BCK(pP) algebra* is a pseudo-BCK algebra \mathfrak{X} satisfying (pP) condition:

(pP) There exists $x \odot y = \min\{z : x \preceq y \rightarrow z\} = \min\{z : y \preceq x \rightsquigarrow z\}$ for all $x, y \in X$.

Definition 2.4 ([16, 20]).

(1) A *pseudo-BCK lattice* is a pseudo-BCK algebra \mathfrak{X} such that $(X; \preceq)$ is a lattice.

(2) A *pseudo-BCK join-semilattice* is a pseudo-BCK algebra \mathfrak{X} such that $(X; \vee)$ is a *join-semilattice*, and $x \rightarrow y = 1$ if and only if $x \vee y = y$.

(3) A *pseudo-BCK meet-semilattice* is a pseudo-BCK algebra \mathfrak{X} such that $(X; \wedge)$ is a *meet-semilattice*, and $x \rightarrow y = 1$ if and only if $x \wedge y = x$.

Definition 2.5 ([16]). A *pseudo-BCK algebra(pDN)* is a bounded pseudo-BCK algebra $\mathfrak{X} = (X; \preceq, \rightarrow, \rightsquigarrow, 0, 1)$ satisfying the condition:

(pDN) $(x \rightarrow) \rightsquigarrow = (x \rightsquigarrow) \rightarrow = x$, where $x \rightarrow = x \rightarrow 0$ and $x \rightsquigarrow = x \rightsquigarrow 0$ for all $x \in X$.

Definition 2.6 ([12]). A pseudo-BCK algebra \mathfrak{X} is called *commutative* if for all $x, y, z \in X$, it satisfies the following identities:

$$(C_1) \quad (x \rightarrow y) \rightsquigarrow y = (y \rightarrow x) \rightsquigarrow x;$$

$$(C_2) \quad (x \rightsquigarrow y) \rightarrow y = (y \rightsquigarrow x) \rightarrow x.$$

Proposition 2.2 ([6]). *Any commutative pseudo-BCK algebra is a join-semilattice with respect to \preceq .*

Theorem 2.2 ([8]). *Let \mathfrak{X} be a pseudo-BCK(pDN) algebra. The following statements are equivalent:*

- (1) $(X; \preceq)$ is a *meet-semilattice*;
- (2) $(X; \preceq)$ is a *join-semilattice*;
- (3) $(X; \preceq)$ is a *lattice*.

Definition 2.7 ([6]). A subset D of a pseudo-BCK algebra \mathfrak{X} is called a *deductive system* of \mathfrak{X} if it satisfies the following conditions:

$$(DS_1) \quad 1 \in D;$$

$$(DS_2) \quad x \in D \text{ and } x \rightarrow y \in D \text{ imply } y \in D.$$

A subset D of \mathfrak{X} is a deductive system if and only if it satisfies (DS_1) and the condition:

$$(DS_3) \quad x \in D \text{ and } x \rightsquigarrow y \in D \text{ imply } y \in D.$$

We will denote by $\mathfrak{D}S(X)$ the set of all deductive systems of \mathfrak{X} .

Definition 2.8 ([6]). A deductive system D of a pseudo-BCK algebra \mathfrak{X} is called *commutative* if it satisfies the following conditions for all $x, y \in X$:

(CDS₁) $y \rightarrow x \in D$ implies $((x \rightarrow y) \rightsquigarrow y) \rightarrow x \in D$;

(CDS₂) $y \rightsquigarrow x \in D$ implies $((x \rightsquigarrow y) \rightarrow y) \rightsquigarrow x \in D$.

We will denote by $\mathfrak{DS}_c(X)$ the set of all commutative deductive systems of a pseudo-BCK algebra \mathfrak{X} .

Definition 2.9 ([3]). A *directoid* is a groupoid $\mathfrak{G} = (G; \vee)$ satisfying the following identities for all $x, y, z \in G$:

- (D₁) $x \vee x = x$;
- (D₂) $(x \vee y) \vee x = x \vee y$;
- (D₃) $y \vee (x \vee y) = x \vee y$;
- (D₄) $x \vee ((x \vee y) \vee z) = (x \vee y) \vee z$.

The relation \leq given by $x \leq y$ if and only if $x \vee y = y$ is a partial order. The binary operation \vee assigns to a pair $\{x, y\}$ is a common upper bound of them.

3. PSEUDO-BCK ALGEBRAS DERIVED FROM DIRECTOIDS

Following the idea used by I. Chajda and J. Kühr [3] for BCK-algebras in what follows we give a generalization of this results for pseudo-BCK algebras. In this section, we assign a semilattice-like structure the sections of which have certain antitone mappings, and also conversely. We have the following results.

Let \mathfrak{X} be a pseudo-BCK algebra. Define binary operations \vee_1 and \vee_2 by:

(A) $x \vee_1 y := (x \rightarrow y) \rightsquigarrow y$ and $x \vee_2 y := (x \rightsquigarrow y) \rightarrow y$ for all $x, y \in X$.

The following examples shows that these operations \vee_1 and \vee_2 need not coincide in general.

Example 3.1 ([8]). Consider the set $X = \{0, a, b, c, 1\}$, where $0 < a, b < c < 1$, a, b incomparable and the operations \rightarrow and \rightsquigarrow given by the following tables:

\rightarrow	0	a	b	c	1		\rightsquigarrow	0	a	b	c	1
0	1	1	1	1	1		0	1	1	1	1	1
a	0	1	b	1	1		a	b	1	b	1	1
b	a	a	1	1	1		b	0	a	1	1	1
c	0	a	b	1	1		c	0	a	b	1	1
1	0	a	b	c	1		1	0	a	b	c	1

Then $\mathfrak{X} = (X; \rightarrow, \rightsquigarrow, 0, 1)$ is a bounded pseudo-BCK algebra, but

$$a \vee_1 0 = (a \rightarrow 0) \rightsquigarrow 0 = 0 \rightsquigarrow 0 = 1 \neq a \vee_2 0 = (a \rightsquigarrow 0) \rightarrow 0 = b \rightarrow 0 = a.$$

Theorem 3.1. *Let \mathfrak{X} be a pseudo-BCK algebra. For every $a \in X$, define unary operations \rightarrow^a and \rightsquigarrow^a by $x \rightarrow^a = x \rightarrow a$ and $x \rightsquigarrow^a = x \rightsquigarrow a$. Then the algebraic structure $\mathfrak{S}(\mathfrak{X}) = (X; \vee_1, \vee_2, (\rightarrow^a)_{a \in X}, (\rightsquigarrow^a)_{a \in X}, 1)$ satisfies the following quasi-identities:*

- (1) $x \vee_1 1 = 1$ and $x \vee_2 1 = 1$;
- (2) $x \rightarrow^x = 1$ and $x \rightsquigarrow^x = 1$;
- (3) $1 \rightarrow^x = x$ and $1 \rightsquigarrow^x = x$;
- (4) $x \vee_1 x = x$ and $x \vee_2 x = x$;

- (5) $(x \vee_1 y) \rightarrow y = x \rightarrow y$ and $(x \vee_2 y) \rightsquigarrow y = x \rightsquigarrow y$;
(6) $x \preceq y$ if and only if $x \vee_1 y = y$ and $x \vee_2 y = y$;
(7) $x \vee_1 y = y$ and $y \vee_1 x = x$ imply $x = y$ and
 $x \vee_2 y = y$ and $y \vee_2 x = x$ imply $x = y$;
(8) $x \vee_1 y = (x \vee_1 y) \vee_1 y = x \vee_1 (x \vee_1 y) = y \vee_1 (x \vee_1 y)$ and
 $x \vee_2 y = (x \vee_2 y) \vee_2 y = x \vee_2 (x \vee_2 y) = y \vee_2 (x \vee_2 y)$;
(9) $(x \vee_1 z) \vee_1 ((x \vee_1 y) \vee_1 z) = (x \vee_1 y) \vee_1 z$ and
 $(x \vee_2 z) \vee_2 ((x \vee_2 y) \vee_2 z) = (x \vee_2 y) \vee_2 z$;
(10) $x \vee_1 y = (x \vee_1 y) \rightarrow y \rightsquigarrow y = ((x \vee_1 y) \rightarrow y \vee_2 y) \rightsquigarrow y$ and
 $x \vee_2 y = (x \vee_2 y) \rightsquigarrow y \rightarrow y = ((x \vee_2 y) \rightsquigarrow y \vee_1 y) \rightarrow y$;
(11) $(x \vee_1 (y \vee_2 z) \rightsquigarrow z) \rightarrow (y \vee_2 z) \rightsquigarrow z = (y \vee_2 (x \vee_1 z) \rightarrow z) \rightsquigarrow (x \vee_1 z) \rightarrow z$;
(12) $(x \vee_1 y) \rightarrow y \vee_2 ((x \vee_1 z) \vee_1 (y \vee_1 z)) \rightarrow (y \vee_1 z) = ((x \vee_1 z) \vee_1 (y \vee_1 z)) \rightarrow (y \vee_1 z)$ and
 $(x \vee_2 y) \rightsquigarrow y \vee_1 ((x \vee_2 z) \vee_2 (y \vee_2 z)) \rightsquigarrow (y \vee_2 z) = ((x \vee_2 z) \vee_2 (y \vee_2 z)) \rightsquigarrow (y \vee_2 z)$;
(13) $((x \vee_2 z) \rightsquigarrow z \vee_1 (y \vee_1 z)) \rightarrow (y \vee_1 z) = ((y \vee_1 z) \rightarrow z \vee_2 (x \vee_2 z)) \rightsquigarrow (x \vee_2 z)$ and
 $((x \vee_1 z) \rightarrow z \vee_2 (y \vee_2 z)) \rightsquigarrow (y \vee_2 z) = ((y \vee_2 z) \rightsquigarrow z \vee_1 (x \vee_1 z)) \rightarrow (x \vee_1 z)$;
(14) $((x \vee_1 y) \vee_1 x) \rightarrow x = (x \vee_1 y) \rightarrow x$ and $((x \vee_2 y) \vee_2 x) \rightsquigarrow x = (x \vee_2 y) \rightsquigarrow x$.

Proof. The proof of (1)-(6) is straightforward by the definition and properties of pseudo-BCK algebras. (6) Assume that $x \preceq y$. Then $x \vee_1 y = (x \rightarrow y) \rightsquigarrow y = 1 \rightsquigarrow y = y$. Also,

$$x \vee_2 y = (x \rightsquigarrow y) \rightarrow y = 1 \rightarrow y = y.$$

Conversely, suppose that $x \vee_1 y = y$ and $x \vee_2 y = y$. Since

$$x \rightarrow y = x \rightarrow ((x \rightarrow y) \rightsquigarrow y) = (x \rightarrow y) \rightsquigarrow (x \rightarrow y) = 1$$

and

$$x \rightsquigarrow y = x \rightsquigarrow ((x \rightsquigarrow y) \rightarrow y) = (x \rightsquigarrow y) \rightarrow (x \rightsquigarrow y) = 1,$$

then $x \preceq y$.

(7) Suppose that $(x \rightarrow y) \rightsquigarrow y = y$ and $(y \rightarrow x) \rightsquigarrow x = x$. Then by (psBCI₃) we have

$$x \rightarrow y = x \rightarrow ((x \rightarrow y) \rightsquigarrow y) = (x \rightarrow y) \rightsquigarrow (x \rightarrow y) = 1$$

and

$$y \rightsquigarrow x = y \rightsquigarrow ((y \rightsquigarrow x) \rightarrow x) = (y \rightsquigarrow x) \rightarrow (y \rightsquigarrow x) = 1.$$

Now, using (psBCI₅) $x = y$. By a similar argument the second part is valid.

(8) By Proposition 2.1 (6), we have

$$\begin{aligned} (x \vee_1 y) \vee_1 y &= ((x \vee_1 y) \rightarrow y) \rightsquigarrow y \\ &= (((x \rightarrow y) \rightsquigarrow y) \rightarrow y) \rightsquigarrow y \\ &= (x \rightarrow y) \rightsquigarrow y \\ &= x \vee_1 y. \end{aligned}$$

Similarly, we see that $x \vee_1 (x \vee_1 y) = x \vee_1 y$ and $y \vee_1 (x \vee_1 y) = x \vee_1 y$.

(9) According to (psBCI₂), $x \preceq x \vee_1 y$. By Proposition 2.1 (1) and (2), we have $x \vee_1 z \preceq (x \vee_1 y) \vee_1 z$. Now, using (6) it follows that $(x \vee_1 z) \vee_1 ((x \vee_1 y) \vee_1 z) = (x \vee_1 y) \vee_1 z$. By a similar argument we can verify $x \vee_2 y = (x \vee_2 y) \vee_2 y = x \vee_2 (x \vee_2 y) = y \vee_2 (x \vee_2 y)$.

(10) $(x \vee_1 y)^{\rightarrow y \rightsquigarrow y} = ((x \vee_1 y) \rightarrow y) \rightsquigarrow y = x \vee_1 y$ and $(x \vee_2 y)^{\rightsquigarrow y \rightarrow y} = ((x \vee_2 y) \rightsquigarrow y) \rightarrow y = x \vee_2 y$. By Proposition 2.1 (6), we have

$$\begin{aligned}
((x \vee_1 y)^{\rightarrow y} \vee_2 y)^{\rightsquigarrow y} &= (((x \vee_1 y) \rightarrow y) \vee_2 y)^{\rightsquigarrow y} \\
&= (((x \rightarrow y) \rightsquigarrow y) \rightarrow y) \vee_2 y)^{\rightsquigarrow y} \\
&= ((x \rightarrow y) \vee_2 y)^{\rightsquigarrow y} \\
&= (((x \rightarrow y) \rightsquigarrow y) \rightarrow y)^{\rightsquigarrow y} \\
&= (x \rightarrow y)^{\rightsquigarrow y} \\
&= (x \rightarrow y) \rightsquigarrow y \\
&= x \vee_1 y.
\end{aligned}$$

Also, the proof of the second part is similar.

(11) From (5) and Proposition 2.1 (4), we conclude

$$\begin{aligned}
(x \vee_1 (y \vee_2 z)^{\rightsquigarrow z})^{\rightarrow (y \vee_2 z)^{\rightsquigarrow z}} &= (x \vee_1 (y \rightsquigarrow z))^{\rightarrow (y \rightsquigarrow z)} \\
&= x \rightarrow (y \rightsquigarrow z) = y \rightsquigarrow (x \rightarrow z) \\
&= (y \vee_2 (x \rightarrow z))^{\rightsquigarrow (x \rightarrow z)} \\
&= (y \vee_2 (x \vee_1 z)^{\rightarrow z})^{\rightsquigarrow (x \vee_1 z)^{\rightarrow z}}.
\end{aligned}$$

(12) Using (5), we have

$$\begin{aligned}
((x \vee_1 z) \vee_1 (y \vee_1 z))^{\rightarrow (y \vee_1 z)} &= (x \vee_1 z) \rightarrow (y \vee_1 z) \\
&= ((x \rightarrow z) \rightsquigarrow z) \rightarrow ((y \rightarrow z) \rightsquigarrow z) \\
&= (y \rightarrow z) \rightsquigarrow (((x \rightarrow z) \rightsquigarrow z) \rightarrow z) \\
&= (y \rightarrow z) \rightsquigarrow (x \rightarrow z).
\end{aligned}$$

We have $(x \rightarrow y) \vee_2 ((y \rightarrow z) \rightsquigarrow (x \rightarrow z)) = 1 \rightarrow ((y \rightarrow z) \rightsquigarrow (x \rightarrow z))$. From this and (psBCI₁) we conclude

$$\begin{aligned}
(x \vee_1 y)^{\rightarrow y} \vee_2 ((x \vee_1 z) \vee_1 (y \vee_1 z))^{\rightarrow (y \vee_1 z)} &= (x \vee_1 y)^{\rightarrow y} \vee_2 ((y \rightarrow z) \rightsquigarrow (x \rightarrow z)) \\
&= (x \rightarrow y) \vee_2 ((y \rightarrow z) \rightsquigarrow (x \rightarrow z)) \\
&= 1 \rightarrow ((y \rightarrow z) \rightsquigarrow (x \rightarrow z)) \\
&= (y \rightarrow z) \rightsquigarrow (x \rightarrow z) \\
&= ((x \vee_1 z) \vee_1 (y \vee_1 z))^{\rightarrow (y \vee_1 z)}.
\end{aligned}$$

(13) Applying (5), we have

$$\begin{aligned}
((x \vee_2 z) \rightsquigarrow^z \vee_1 (y \vee_1 z)) \rightarrow^{(y \vee_1 z)} &= (x \rightsquigarrow z) \rightarrow (y \vee_1 z) \\
&= (x \rightsquigarrow z) \rightarrow ((y \rightarrow z) \rightsquigarrow z) \\
&= (y \rightarrow z) \rightsquigarrow ((x \rightsquigarrow z) \rightarrow z) \\
&= (y \rightarrow z) \rightsquigarrow (x \vee_2 z) \\
&= ((y \vee_1 z) \rightarrow^z \vee_2 (x \vee_2 z)) \rightsquigarrow^{(x \vee_2 z)}.
\end{aligned}$$

By a similar argument we have

$$((x \vee_1 z) \rightarrow^z \vee_2 (y \vee_2 z)) \rightsquigarrow^{(y \vee_2 z)} = ((y \vee_2 z) \rightsquigarrow^z \vee_1 (x \vee_1 z)) \rightarrow^{(x \vee_1 z)}.$$

(14) Using Proposition 2.1 (6), we get

$$\begin{aligned}
((x \vee_1 y) \vee_1 x) \rightarrow^x &= (((x \rightarrow y) \rightsquigarrow y) \vee_1 x) \rightarrow^x \\
&= (((x \rightarrow y) \rightsquigarrow y) \rightarrow x) \rightsquigarrow x \rightarrow^x \\
&= ((x \rightarrow y) \rightsquigarrow y) \rightarrow^x \\
&= ((x \rightarrow y) \rightsquigarrow y) \rightarrow x \\
&= (x \vee_1 y) \rightarrow x \\
&= (x \vee_1 y) \rightarrow^x.
\end{aligned}$$

Similarly, $((x \vee_2 y) \vee_2 x) \rightsquigarrow^x = (x \vee_2 y) \rightsquigarrow^x$. \square

Lemma 3.1. *Let $\mathfrak{X} = (X; \vee_1, \vee_2)$ be an algebra of type $(2, 2)$ satisfying the quasi-identities (4), (7), (8) and (9) of Theorem 3.1. Then the binary relation \preceq defined by*

(B) $x \preceq y$ if and only if $x \vee_1 y = y$ and $x \vee_2 y = y$ is a partial order on X .

Proof. By (4) and (7), \preceq is reflexive and antisymmetric. For transitivity, assume that $x \preceq y$ and $y \preceq z$. Using (8) and (9), we get

$$\begin{aligned}
x \vee_1 z &= (x \vee_1 z) \vee_1 z \\
&= (x \vee_1 z) \vee_1 (y \vee_1 z) \\
&= (x \vee_1 z) \vee_1 ((x \vee_1 y) \vee_1 z) \\
&= (x \vee_1 y) \vee_1 z \\
&= y \vee_1 z = z
\end{aligned}$$

and if $x \vee_2 y = y$ and $y \vee_2 z = z$, then we have

$$\begin{aligned}
x \vee_2 z &= (x \vee_2 z) \vee_2 z \\
&= (x \vee_2 z) \vee_2 (y \vee_2 z) \\
&= (x \vee_2 z) \vee_2 ((x \vee_2 y) \vee_2 z) \\
&= (x \vee_2 y) \vee_2 z \\
&= y \vee_2 z = z.
\end{aligned}$$

Thus, \preceq is a partial order on X . □

The following example shows that for every pseudo-BCK algebra \mathfrak{X} , $(X; \vee_1)$ and $(X; \vee_2)$ are not directoids in general.

Example 3.2. Let \mathfrak{X} be the algebra given in Example 3.1. Then $(X; \vee_1)$ and $(X; \vee_2)$ are not directoids, since

$$c \vee_1 0 = (c \rightarrow 0) \rightsquigarrow 0 = 0 \rightsquigarrow 0 = 1 \neq (c \vee_1 0) \vee_1 c = 0 \vee_1 c = (0 \rightarrow c) \rightsquigarrow c = 1 \rightsquigarrow c = c$$

and

$$c \vee_2 0 = (c \rightsquigarrow 0) \rightarrow 0 = 0 \rightarrow 0 = 1 \neq (c \vee_2 0) \vee_2 c = 0 \vee_2 c = (0 \rightsquigarrow c) \rightarrow c = 1 \rightarrow c = c.$$

Theorem 3.2. *Let \mathfrak{X} be a pseudo-BCK algebra, \vee_1 and \vee_2 be the binary operations defined by (A). Then the following conditions are equivalent:*

- (1) $(X; \vee_1)$ and $(X; \vee_2)$ are directoids;
- (2) \mathfrak{X} is a commutative pseudo-BCK algebra;
- (3) $(X; \preceq)$ is a join-semilattice, where \preceq is defined by (B).

Proof. (1) \Rightarrow (2) Assume that $(X; \vee_1)$ is a directoid. Then $x \preceq y$ implies $y \vee_1 x = y$ and so \mathfrak{X} satisfies the quasi-identity

$$x \preceq y \Rightarrow y = (y \rightarrow x) \rightsquigarrow x.$$

Similarly, $x \preceq y$ implies $y = (y \rightsquigarrow x) \rightarrow x$. Therefore, \mathfrak{X} is a commutative pseudo-BCK algebra by Theorem 2.1.

(2) \Rightarrow (3) It follows from Proposition 2.2.

(3) \Rightarrow (1) It is obvious that every join-semilattice is a directoid. □

Corollary 3.1. *Let \mathfrak{X} be a pseudo-BCK(pDN), \vee_1 and \vee_2 be the binary operations defined by (A). Then the following conditions are equivalent:*

- (1) $(X; \vee_1)$ and $(X; \vee_2)$ are directoids;
- (2) \mathfrak{X} is a commutative pseudo-BCK algebra;
- (3) $(X; \preceq)$ is a join-semilattice;
- (4) $(X; \preceq)$ is a meet-semilattice;
- (5) $(X; \preceq)$ is a lattice.

Proof. It follows from Theorems 3.2 and 2.2. □

Corollary 3.2. *Let \mathfrak{X} be a pseudo-BCK(pDN), \vee_1 and \vee_2 be the binary operations defined by (A). Then the following conditions are equivalent:*

- (1) $(X; \vee_1)$ and $(X; \vee_2)$ are directoids;
- (2) \mathfrak{X} is a commutative pseudo-BCK algebra;
- (3) $\{1\} \in \mathfrak{D}S_c(X)$;
- (4) $\mathfrak{D}S(X) = \mathfrak{D}S_c(X)$.

Proof. It follows from Theorem 3.2 and [6, Corollary 4.6, Theorem 4.7 and Corollary 4.8]. □

In [8], L. C. Ciungu proved that for every pseudo-BCK(pDN) lattice the following conditions are equivalent (see [8, Proposition 3.5]):

- (P₁) $(x \wedge y) \rightarrow z = (x \rightarrow z) \vee (y \rightarrow z)$ and $(x \wedge y) \rightsquigarrow z = (x \rightsquigarrow z) \vee (y \rightsquigarrow z)$;
(P₂) $z \rightarrow (x \vee y) = (z \rightarrow x) \vee (z \rightarrow y)$ and $z \rightsquigarrow (x \vee y) = (z \rightsquigarrow x) \vee (z \rightsquigarrow y)$.

Also, she showed that the class of pseudo-BCK(pDN) lattices satisfies the conditions (P₁) and (P₂) is not empty, since every pseudo-MV algebra satisfies these conditions. Further, It was proved that if a pseudo-BCK(pDN) lattice \mathfrak{X} satisfying (P₁) or (P₂), then $(X; \preceq)$ is a distributive lattice (see [8, Theorem 3.4, Corollary 3.2]).

Theorem 3.3. *Let $\mathfrak{S} = (S; \vee_1, \vee_2, (\rightarrow^a)_{a \in S}, (\rightsquigarrow^a)_{a \in S}, 1)$ be a structure algebraic, where \vee_1 and \vee_2 are binary operations on S and for each $a \in S$, \rightarrow^a and \rightsquigarrow^a are unary operations on $\{x \in S : a \vee_1 x = 1 \text{ and } a \vee_2 x = 1\}$ and 1 is a distinguished element of S , satisfying the quasi-identities (1)-(12) from Theorem 3.1. Define the new binary operations \rightarrow and \rightsquigarrow on S by*

$$(C) \quad x \rightarrow y = (x \vee_1 y)^{\rightarrow y} \text{ and } x \rightsquigarrow y = (x \vee_2 y)^{\rightsquigarrow y}.$$

Then $\mathfrak{X}(\mathfrak{S}) = (S; \rightarrow, \rightsquigarrow, 1)$ is a pseudo-BCK algebra.

Proof. The definition of \rightarrow and \rightsquigarrow are well define from Theorem 3.1 (8). Furthermore, we note that

$$(D) \quad x \vee_1 y = y \text{ and } x \vee_2 y = y \text{ if and only if } x \rightarrow y = 1 \text{ and } x \rightsquigarrow y = 1.$$

Indeed, if $x \vee_1 y = y$, then $x \rightarrow y = (x \vee_1 y)^{\rightarrow y} = y^{\rightarrow y} = 1$, by Theorem 3.1 (2). Similarly, if $x \vee_2 y = y$, then $x \rightsquigarrow y = (x \vee_2 y)^{\rightsquigarrow y} = y^{\rightsquigarrow y} = 1$.

For conversely, $1 = x \rightarrow y = (x \vee_1 y)^{\rightarrow y}$ implies $y = 1^{\rightsquigarrow y} = (x \vee_1 y)^{\rightarrow y \rightsquigarrow y} = x \vee_1 y$. Also, $1 = x \rightsquigarrow y = (x \vee_2 y)^{\rightsquigarrow y}$ implies $y = 1^{\rightarrow y} = (x \vee_2 y)^{\rightsquigarrow y \rightarrow y} = x \vee_2 y$, by Theorem 3.1 (3) and (10). Now, we verify the axioms of pseudo-BCK algebras as follows.

(psBCI₁) Using Theorem 3.1 (5) and (12), we obtain

$$\begin{aligned} (x \rightarrow y) \vee_2 ((x \vee_1 z) \rightarrow (y \vee_1 z)) &= (x \vee_1 y)^{\rightarrow y} \vee_2 ((x \vee_1 z) \vee_1 (y \vee_1 z))^{\rightarrow (y \vee_1 z)} \\ &= ((x \vee_1 z) \vee_1 (y \vee_1 z))^{\rightarrow (y \vee_1 z)} \\ &= (x \vee_1 z) \rightarrow (y \vee_1 z). \end{aligned}$$

Thus, $(x \rightarrow y) \rightsquigarrow ((x \vee_1 z) \rightarrow (y \vee_1 z)) = 1$.

Also, according to Theorem 3.1 (10) and (11), we get

$$(x \rightarrow y) \rightsquigarrow y = ((x \vee_1 y)^{\rightarrow y} \vee_2 y)^{\rightsquigarrow y} = x \vee_1 y$$

and

$$x \rightarrow (y \rightsquigarrow z) = (x \vee_1 (y \vee_2 z)^{\rightsquigarrow z})^{\rightarrow (y \vee_2 z)} = (y \vee_2 (x \vee_1 z)^{\rightarrow z})^{\rightsquigarrow (x \vee_1 z)} = y \rightsquigarrow (x \rightarrow z).$$

Then

$$\begin{aligned} (x \vee_1 z) \rightarrow (y \vee_1 z) &= ((x \rightarrow z) \rightsquigarrow z) \rightarrow ((y \rightarrow z) \rightsquigarrow z) \\ &= (y \rightarrow z) \rightsquigarrow (((x \rightarrow z) \rightsquigarrow z) \rightarrow z) \\ &= (y \rightarrow z) \rightsquigarrow (x \rightarrow z), \end{aligned}$$

since

$$\begin{aligned} ((x \rightarrow z) \rightarrow z) \rightsquigarrow z &= (((x \vee_1 z)^{\rightarrow z} \vee_2 z)^{\rightarrow z} \vee_1 z)^{\rightarrow z} \\ &= ((x \vee_1 z) \vee_1 z)^{\rightarrow z} \\ &= (x \vee_1 z)^{\rightarrow z} \\ &= x \rightarrow z. \end{aligned}$$

Altogether, we have

$$(x \rightarrow y) \rightsquigarrow ((y \rightarrow z) \rightsquigarrow (x \rightarrow z)) = (x \rightarrow y) \rightsquigarrow ((x \vee_1 z) \rightarrow (y \vee_1 z)) = 1.$$

The second part of axiom (psBCI₁) follows by duality.

(psBCI₂) Using Theorem 3.1 (10), we get

$$(x \rightarrow y) \rightsquigarrow y = ((x \vee_1 y)^{\rightarrow y} \vee_2 y)^{\rightsquigarrow y} = x \vee_1 y.$$

Hence, $x \vee_1 ((x \rightarrow y) \rightsquigarrow y) = x \vee_1 (x \vee_1 y) = x \vee_1 y = (x \rightarrow y) \rightsquigarrow y$. Then $x \rightarrow ((x \rightarrow y) \rightsquigarrow y) = 1$.

By a similar argument we have $x \rightsquigarrow ((x \rightsquigarrow y) \rightarrow y) = 1$.

(psBCI₃) Applying (D), from $x \vee_1 x = x \vee_2 x = x$ it follows $x \rightarrow x = x \rightsquigarrow x = 1$.

(psBCI₄) If $x \rightarrow y = 1$ and $y \rightarrow x = 1$, then by Theorem 3.1 (6), we have $x \vee_1 y = y$ and $y \vee_1 x = x$. Now, using Theorem 3.1 (7), it follows $x = y$.

(psBCI₅) This follows from Theorem 3.1 (6) and (D).

(psBCK) By Theorem 3.1 (1), $x \vee_1 1 = x \vee_2 1 = 1$. From (D) we see that $x \rightarrow 1 = x \rightsquigarrow 1 = 1$. \square

Theorem 3.4. *Let $\mathfrak{S} = (S; \vee_1, \vee_2, (\rightarrow^a)_{a \in S}, (\rightsquigarrow^a)_{a \in S}, 1)$ be an algebra as in Theorem 3.3 satisfying (1)-(12) of Theorem 3.1 and \mathfrak{X} be a pseudo-BCK algebra. Then $\mathfrak{X}(\mathfrak{S}(\mathfrak{X})) = \mathfrak{X}$ and $\mathfrak{S}(\mathfrak{X}(\mathfrak{S})) = \mathfrak{S}$.*

Proof. By Theorem 3.1, $\mathfrak{S}(\mathfrak{X}) = (X; \vee_1, \vee_2, (\rightarrow^a)_{a \in X}, (\rightsquigarrow^a)_{a \in X}, 1)$ is the structure satisfying (1)-(12) which is assigned to a given pseudo-BCK algebra \mathfrak{X} . Then in $\mathfrak{X}(\mathfrak{S}(\mathfrak{X})) = (X; \rightarrow_1, \rightsquigarrow_1, 1)$ we have

$$x \rightarrow_1 y = (x \vee_1 y)^{\rightarrow y} = ((x \rightarrow y) \rightsquigarrow y) \rightarrow y = x \rightarrow y$$

and

$$x \rightsquigarrow_2 y = (x \vee_2 y)^{\rightsquigarrow y} = ((x \rightsquigarrow y) \rightarrow y) \rightsquigarrow y = x \rightsquigarrow y.$$

Therefore, $\mathfrak{X}(\mathfrak{S}(\mathfrak{X})) = \mathfrak{X}$.

Also, assume that $\mathfrak{S} = (S; \vee_1, \vee_2, (\rightarrow^a)_{a \in S}, (\rightsquigarrow^a)_{a \in S}, 1)$ is a structure that satisfies (1)-(12) of Theorem 3.1, $\mathfrak{X}(\mathfrak{S}) = (S; \rightarrow, \rightsquigarrow, 1)$ its corresponding pseudo-BCK algebra (cf. Theorem 3.3) and $\mathfrak{S}(\mathfrak{X}(\mathfrak{S})) = (S; \sqcup_1, \sqcup_2, (r_{1a})_{a \in S}, (r_{2a})_{a \in S}, 1)$. Then

$$x \sqcup_1 y = (x \rightarrow y) \rightsquigarrow y = ((x \vee_1 y)^{\rightarrow y} \vee_2 y)^{\rightsquigarrow y} = x \vee_1 y$$

and

$$x \sqcup_2 y = (x \rightsquigarrow y) \rightarrow y = ((x \vee_2 y)^{\rightsquigarrow y} \vee_1 y)^{\rightarrow y} = x \vee_2 y.$$

Further, for $x \in [a, 1]$, we have

$$r_{1a}(x) = x \rightarrow a = (x \vee_1 a)^{\rightarrow a} = ((a \vee_1 x) \vee_1 a)^{\rightarrow a} = (a \vee_1 x)^{\rightarrow a} = x^{\rightarrow a}$$

and

$$r_{2a}(x) = x \rightsquigarrow a = (x \vee_2 a)^{\rightsquigarrow a} = ((a \vee_2 x) \vee_2 a)^{\rightsquigarrow a} = (a \vee_2 x)^{\rightsquigarrow a} = x^{\rightsquigarrow a}.$$

Therefore, $\mathfrak{S}(\mathfrak{X}(\mathfrak{S})) = \mathfrak{S}$. □

Corollary 3.3. *Let $\mathfrak{S} = (S; \vee_1, \vee_2, (\rightarrow^a)_{a \in S}, (\rightsquigarrow^a)_{a \in S}, 1)$ be an algebraic structure satisfying (1)-(13) of Theorem 3.1. Then the relation defined by (B) is a partial order on S , 1 is the greatest element of S and for every $x, y \in S$, $x, y \preceq x \vee y$, where $\vee = \vee_1 = \vee_2$. Moreover, for each $a \in S$, \rightarrow^a and \rightsquigarrow^a are antitone mappings on $[a, 1] = \{x \in S : a \preceq x\}$.*

CONCLUSION

We consider that this paper could contribute to the study of algebraic structures and to the development of pseudo-BCK algebras. So, we hope it would be served as a foundation and another topic of research to define and investigate among algebraic structures derived from pseudo-BCK algebras. As another direction of research, one could investigate relationship between commutative pseudo-valuation on pseudo-BCK algebras with directoids.

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