

Bulova logika

$$T = (\lambda x . \lambda y . x)$$

$$F = (\lambda x . \lambda y . y)$$

- Nije obavezno da tačno i netačno budu definisani na ovaj način

$$\mathit{and} = (\lambda a . \lambda b . a b F)$$

- $\mathit{and} T T$
- $(\lambda a . \lambda b . a b (\lambda x . \lambda y . y)) T T$

and T T

$(\lambda a. \lambda b. a b (\lambda x. \lambda y. y)) (\lambda x. \lambda y. x) (\lambda x. \lambda y. x)$

$(\lambda b. a b (\lambda x. \lambda y. y)) [a \rightarrow (\lambda x. \lambda y. x)] (\lambda x. \lambda y. x)$

$(\lambda b. (\lambda x. \lambda y. x) b (\lambda x. \lambda y. y)) (\lambda x. \lambda y. x)$

$((\lambda x. \lambda y. x) b (\lambda x. \lambda y. y)) [b \rightarrow (\lambda x. \lambda y. x)]$

$(\lambda x. \lambda y. x)(\lambda x. \lambda y. x) (\lambda x. \lambda y. y)$

$(\lambda y. x)[x \rightarrow (\lambda x. \lambda y. x)] (\lambda x. \lambda y. y)$

$(\lambda y. (\lambda x. \lambda y. x)) (\lambda x. \lambda y. y)$

$(\lambda x. \lambda y. x) [y \rightarrow (\lambda x. \lambda y. y)]$

$(\lambda x. \lambda y. x)$

T

and T F

$(\lambda a. \lambda b. a b F) T F$

$(\lambda b. a b F) [a \rightarrow T] F$

$(\lambda b. T b F) F$

$(T b F)[b \rightarrow F]$

$(T F F)$

$(\lambda x. \lambda y. x) F F$

$(\lambda y. x) [x \rightarrow F] F$

$(\lambda y. F) F$

$F [y \rightarrow F]$

F

and F T

$(\lambda a. \lambda b. a b F) F T$

$(F T F)$

F

and F F

$(\lambda a. \lambda b. a b F) F F$

$(F F F)$

F

not

- $\text{not } T = F$
- $\text{not } F = T$

$\text{not} = (\lambda a . a F T)$

$\text{not } T$

- $(\lambda a . a F T) T$
- $T F T$
- F

$\text{not } F$

- $(\lambda a . a F T) F$
- $F F T$
- T

IF grananje

```
if c then
  a
else
  b
```

if c a b

if T a b = a

if F a b = b

if = ($\lambda a . a$)

Primeri

if T a b

- $(\lambda a . a) T a b$
- $T a b$
- a

if F a b

- $(\lambda a . a) F a b$
- $F a b$
- b

Čerčovi numerali

$$0 = \lambda f . \lambda x . x$$

$$1 = \lambda f . \lambda x . f x$$

$$2 = \lambda f . \lambda x . f f x$$

$$3 = \lambda f . \lambda x . f f f x$$

$$4 = \lambda f . \lambda x . f f f f x$$

- $\lambda f . \lambda x . (f (f (f (f x))))$

$$4 a b$$

- $a a a a b$

Funkcija naslednika

$succ = \lambda n . \lambda f . \lambda x . f (n f x)$

$0 = \lambda f . \lambda x . x$

$succ\ 0$

- $\lambda n . \lambda f . \lambda x . f (n f x)\ 0$
- $\lambda f . \lambda x . f (0 f x)$
- $\lambda f . \lambda x . f ((\lambda f . \lambda x . x) f x)$
- $\lambda f . \lambda x . f x$
- 1

Funkcija naslednika

$$\mathit{succ} = \lambda n . \lambda f . \lambda x . f (n f x)$$

$$1 = \lambda f . \lambda x . f x$$

$$\mathit{succ} 1$$

- $\lambda n . \lambda f . \lambda x . f (n f x) 1$
- $\lambda f . \lambda x . f (1 f x)$
- $\lambda f . \lambda x . f ((\lambda f . \lambda x . f x) f x)$
- $\lambda f . \lambda x . f f x$

$$2 = \lambda f . \lambda x . f f x$$

$$\mathit{succ} 1 = 2$$

$$\mathit{succ} n = n + 1$$

Sabiranje

$add\ 0\ 1 = 1$

$add\ 1\ 2 = 3$

$add = \lambda n . \lambda m . \lambda f . \lambda x . n\ f\ (m\ f\ x)$

$add\ 0\ 1$

- $(\lambda n . \lambda m . \lambda f . \lambda x . n\ f\ (m\ f\ x))\ 0\ 1$
- $(\lambda m . \lambda f . \lambda x . 0\ f\ (m\ f\ x))\ 1$
- $\lambda f . \lambda x . 0\ f\ (1\ f\ x)$
- $\lambda f . \lambda x . 0\ f\ (f\ x)$
- $\lambda f . \lambda x . f\ x$

Množenje

mult 0 1 = 0

mult 1 2 = 2

mult 2 5 = 10

mult = $\lambda n . \lambda m . m (add\ n) 0$

mult 0 1

- $(\lambda n . \lambda m . m (add\ n) 0) 0 1$
- $(\lambda m . m (add\ 0) 0) 1$
- $1 (add\ 0) 0$
- $add\ 0\ 0$
- 0

Množenje

mult 1 2

- $(\lambda n . \lambda m . m (add\ n)\ 0)\ 1\ 2$
- $(\lambda m . m (add\ 1)\ 0)\ 2$
- $2 (add\ 0)\ 0$
- $(add\ 1)\ ((add\ 1)\ 0)$
- $(add\ 1)\ (add\ 1\ 0)$
- $(add\ 1)\ (1)$

Faktorijel

$n!$

- $\text{fact}(0) = 1$
- $\text{fact}(n) = n * \text{fakt}(n-1)$

$\text{fact} = \left(\lambda n . \text{if } (\text{iszero } n) (1) \left(\text{mult } n \left(\text{fact } (\text{pred } n) \right) \right) \right)$

- Funkciju ne možemo definisati ovako!

```
int fact(int n)
{
    if (n==0)
        return 1;
    return n * fact(n-1);
}
```

Y kombinator

$$Y = (\lambda x . \lambda y . y (x x y)) (\lambda x . \lambda y . y (x x y))$$

Y foo

- $(\lambda x . \lambda y . y (x x y)) (\lambda x . \lambda y . y (x x y)) \text{foo}$
- $(\lambda y . y ((\lambda x . \lambda y . y (x x y)) (\lambda x . \lambda y . y (x x y)) y)) \text{foo}$
- $\text{foo} ((\lambda x . \lambda y . y (x x y)) (\lambda x . \lambda y . y (x x y)) \text{foo})$
- $\text{foo} (Y \text{foo})$
- $\text{foo} (\text{foo} (Y \text{foo}))$
- $\text{foo} (\text{foo} (\text{foo} (Y \text{foo})))$
- ...

Rekurzija

$$\times fact = (\lambda n . if (iszero n) (1) (mult n (fact (pred n))))$$

$$fact = Y (\lambda f . \lambda n . if (iszero n) (1) (mult n (f (pred n))))$$

fact 1

- $Y (\lambda f . \lambda n . if (iszero n) (1) (mult n (f (pred n)))) 1$
- $(\lambda f . \lambda n . if (iszero n) (1) (mult n (f (pred n)))) (Y (\lambda f . \lambda n . if (iszero n) (1) (mult n (f (pred n)))) 1)$
- $(\lambda n . if (iszero n) (1) (mult n ((Y (\lambda f . \lambda n . if (iszero n) (1) (mult n (f (pred n)))) (pred n)))) 1$
- $if (iszero 1) (1) (mult 1 ((Y (\lambda f . \lambda n . if (iszero n) (1) (mult n (f (pred n)))) (pred 1))))$
- $if F (1) (mult 1 ((Y (\lambda f . \lambda n . if (iszero n) (1) (mult n (f (pred n)))) (pred 1))))$
- $mult 1 ((Y (\lambda f . \lambda n . if (iszero n) (1) (mult n (f (pred n)))) (pred 1)))$

Rekurzija

fact 1

- $mult\ 1\ \left(\left(Y\ (\lambda\ f.\ \lambda\ n.\ if\ (iszero\ n)\ (1)\ (mult\ n\ (f\ (pred\ n)))) \right) (pred\ 1) \right)$
- $mult\ 1\ \left(\left(\lambda\ f.\ \lambda\ n.\ if\ (iszero\ n)\ (1)\ (mult\ n\ (f\ (pred\ n))) \right) \left(Y\ (\lambda\ f.\ \lambda\ n.\ if\ (iszero\ n)\ (1)\ (mult\ n\ (f\ (pred\ n)))) \right) (pred\ 1) \right)$
- $mult\ 1\ \left(\lambda\ n.\ if\ (iszero\ n)\ (1)\ \left(mult\ n\ \left(\left(Y\ (\lambda\ f.\ \lambda\ n.\ if\ (iszero\ n)\ (1)\ (mult\ n\ (f\ (pred\ n)))) \right) (pred\ n) \right) \right) (pred\ 1) \right)$
- $mult\ 1\ \left(\lambda\ n.\ if\ (iszero\ n)\ (1)\ \left(mult\ n\ \left(\left(Y\ (\lambda\ f.\ \lambda\ n.\ if\ (iszero\ n)\ (1)\ (mult\ n\ (f\ (pred\ n)))) \right) (pred\ n) \right) \right) 0 \right)$
- $mult\ 1\ \left(if\ (iszero\ 0)\ (1)\ \left(mult\ 0\ \left(\left(Y\ (\lambda\ f.\ \lambda\ n.\ if\ (iszero\ n)\ (1)\ (mult\ n\ (f\ (pred\ n)))) \right) (pred\ 0) \right) \right) \right)$
- $mult\ 1\ \left(if\ T\ (1)\ \left(mult\ 0\ \left(\left(Y\ (\lambda\ f.\ \lambda\ n.\ if\ (iszero\ n)\ (1)\ (mult\ n\ (f\ (pred\ n)))) \right) (pred\ 0) \right) \right) \right)$
- $mult\ 1\ 1$
- 1

Tjuring-kompletan sistem

Bulova
logika

Aritmetika

Petlje

Haskell

GRAHAM HUTTON, PROGRAMMING IN HASKELL,
CAMBRIDGE UNIVERSITY PRESS (2007)

Haskell kao kalkulator

```
Prelude> 2+3*4
14
Prelude> (2+3)*4
20
Prelude> sqrt(3^2 + 4^2)
5.0
Prelude> 2^1000
107150860718626732094842504906000181056140481170553360744375038837035105112493612249319837881569585812759467291755314682
518714528569231404359845775746985748039345677748242309854210746050623711418779541821530464749835819412673987675591655439
46077062914571196477686542167660429831652624386837205668069376
Prelude>
```

```
Prelude> (+) 2 3
5
Prelude> (^) 2 5
32
```

Standardne funkcije

Prvi element liste

```
Prelude> head [1, 2, 3, 4]  
1
```

Svi elementi sem prvog

```
Prelude> tail [1, 2, 3, 4]  
[2,3,4]
```

N-ti element liste

```
Prelude> [1, 2, 3, 4] !! 2  
3
```

Prvih n elemenata liste

```
Prelude> take 3 [1, 2, 3, 4]  
[1,2,3]
```

Standardne funkcije

Uklanjanja prvih n elemenata liste

```
Prelude> drop 3 [1, 2, 3, 4]  
[4]
```

```
xs = take n xs ++ drop n xs
```

Dužina liste

```
Prelude> length [1, 2, 3, 4]  
4
```

Suma elemenata liste

```
Prelude> sum [1, 2, 3, 4]  
10
```

Standardne funkcije

Proizvod elemenata liste

```
Prelude> product [1, 2, 3, 4]  
24
```

Nadovezivanje lista

```
Prelude> [1, 2, 3] ++ [4, 5]  
[1,2,3,4,5]
```

Inverzna lista

```
Prelude> reverse [1, 2, 3, 4]  
[4,3,2,1]
```


Funkcijsko vs. objektno programiranje

drop 3 [1, 2, 3, 4, 5] vs. [1, 2, 3, 4, 5].drop(3)

method receiver a b c vs. **receiver**.method(a,b,c)

OO programiranje lakše „teče“ – metode objekta navode na način rešavanja

vs.

U funkcijskom programiranju programer preuzima inicijativu

Svi argumenti su jednako važni

Matematika vs. funkcijsko programiranje

$f(a, b) + c d$ vs. $f a b + c * d$

- Primeniti funkciju f na argumente a i b i dodati proizvod $c d$
- Najčešće korišćen simbol zauzima najmanje mesta -- proizvod vs. aplikacija
- Prioritet operacija -- proizvod vs. aplikacija

```
ghci> let f x = x + 3
ghci> f 5 * 7
56
```

Elegantna sintaksa!

Zagrade mogu dodatno da se izbegnu korišćenjem $\$$ za definisanje desne asocijativnosti

Math	Haskell
$f(x)$	<code>f x</code>
$f(x, y)$	<code>f x y</code>
$f(g(x))$	<code>f (g x)</code>
$f(x, g(y))$	<code>f x (g x)</code>
$f(x)g(x)$	<code>f x * g x</code>

Skript fajl

```
double x = 2 * x
quadruple x = double (double x)
```

```
double x = 2 * x
quadruple = double . double
```

```
ghci> :load Primeri01_1.hs
[1 of 1] Compiling Main                ( Primeri01_1.hs, interpreted )
Ok, one module loaded.
ghci> quadruple 2
8
ghci> :type quadruple
quadruple :: Integer -> Integer
ghci> take (double 3) [1..15]
[1,2,3,4,5,6]
```

Kako u C++/C# izgleda definicija i testiranje ovakvog koda?

Glavni delovi koda i „šum“

Skript fajl

```
faktor n = product [1 .. n]
```

- U proceduralnim jezicima petlja ili rekurzija
- $1*2*...*n$
- `[1 .. n]`
- `[1 ..]`
- `take n [1 ..]`

```
faktor n = product (take n [1 .. ])
```

```
average ns = sum ns `div` length ns
```

- Korišćenje funkcija kao infiksnih operatora

```
average ns = div (sum ns) (length ns)
```

Imenovanje

Nazivi funkcija i argumenata počinju malim slovom

Imena tipova počinju velikim slovom

Konvencija

- n – broj
- xs – lista
- xss – lista lista

Izgled

```
a = 5  
b = 10  
c = 30
```



```
a = 5  
  b = 10  
c = 30
```



```
a = 5  
b = 10  
  c = 30
```



Izbegava se pisanje dodatnih simbola za ograničavanja blokova

Tipovi i klase

Tipovi

Tip je ime za kolekciju povezanih vrednosti

Osnovni tipovi

- **Bool** = False | True

```
Prelude> 1 + False
<interactive>:16:1: error:
  * No instance for (Num Bool) arising from a use of `+'
  * In the expression: 1 + False
    In an equation for `it': it = 1 + False
```

- **Char** = 'a' | 'b' | ... | 'A' | 'B' | ...
- **String**
- **Int** = -2^{31} | ... | -1 | 0 | 1 | ... | $2^{31}-1$
- **Integer** (neograničen tip, proizvoljno velike vrednosti)
- **Float**
- **Double**

Tipovi

exp :: Type

```
ghci> [1,2,3] :: [Int]
[1,2,3]
ghci> :type [1,2,3]
[1,2,3] :: Num a => [a]
ghci> :type [['a'],['b','c']]
[['a'],['b','c']] :: [[Char]]
ghci> :type ('a', True, 1)
('a', True, 1) :: Num c => (Char, Bool, c)
ghci> :type ('a', (False, "abc"))
('a', (False, "abc")) :: (Char, (Bool, String))
ghci> :type [1, 2.2, 3]
[1, 2.2, 3] :: Fractional a => [a]
```

Definisanje tipova nije obavezno

Može donekle da ubrza kod

Tipovi lista

Lista je niz elemenata istog tipa

- `[False, True, False] :: [Bool]`
- `['a', 'b', 'c', 'd'] :: [Char]`

`[t]` je lista elemenata tipa `t`

Ista notacija za konstruktor tipa i konstruktor vrednosti

Tip liste ništa ne govori o dužini liste

`[[[Char]]] ?`

Tipovi torki

Torka je niz vrednosti različitog tipa

- `(False, True) :: (Bool, Bool)`
- `(False, 'a', True) :: (Bool, Char, Bool)`
- `(1, True, 'a') :: (Int, Bool, Char)`

`(t1, t2, ..., tn)` je n-torka čija i-ta komponenta ima tip `ti` za svako `i` u `[1..n]`

Ista notacija za konstruktor tipa i konstruktor vrednosti

Definica tipa definiše i dužinu/veličinu

Ne postoji ograničenje za tip komponente

- `('a', (False, 'b')) :: (Char, (Bool, Char))`
- `(True, ['a', 'b']) :: (Bool, [Char])`