

Monadi i ostalo

Funktori

Uopštenje funkcije map

Ideja mapiranja funkcija na svaki element ne mora biti ograničena na liste

Funktori – klasa tipova koji podržavaju upšteno mapiranja

```
map :: (a -> b) -> [a] -> [b]
map f xs = [f x | x <- xs]
```

```
class Functor f where
  fmap :: (a -> b) -> f a -> f b
```

fmap prihvata funkciju $a \rightarrow b$ i na strukturu tipa $f\ a$, čiji su elementi tipa a , primenjuje funkciju i vraća strukturu tipa $f\ b$

f mora biti parametrizovan tip

Funktori

```
instance Functor [] where
-- fmap :: (a -> b) -> [a] -> [b]
  fmap = map
```

```
instance Functor Maybe where
-- fmap :: (a -> b) -> Maybe a -> Maybe b
  fmap _ Nothing = Nothing
  fmap g (Just x) = Just (g x)
```

```
ghci> map (+1) [1,3,4,7]
[2,4,5,8]
ghci> fmap (+1) [1,3,4,7]
[2,4,5,8]
```

```
ghci> fmap (+1) Nothing
Nothing
ghci> fmap (^2) (Just 3)
Just 9
ghci> fmap not (Just False)
Just True
```

```
ghci> fmap (++"!") getLine
hi
"hi!"
```

Funktori

```
data Tree a = Leaf a | Node (Tree a) (Tree a)
  deriving Show

instance Functor Tree where
  -- fmap :: (a -> b) -> Tree a -> Tree b
  fmap f (Leaf x) = Leaf (f x)
  fmap f (Node left right) = Node (fmap f left) (fmap f right)
```

```
ghci> fmap length (Leaf "abc")
Leaf 3
ghci> fmap even (Node (Leaf 1) (Leaf 2))
Node (Leaf False) (Leaf True)
ghci> fmap (\x -> x*2) (Node (Leaf 1) (Leaf 2))
Node (Leaf 2) (Leaf 4)
```

Zakoni za funktore

`fmap id = id`

`fmap (g.h) = fmap g . fmap h`

```
ghci> fmap (not.even) [1,2]
[True,False]
ghci> (fmap not . fmap even) [1,2]
[True,False]
```

Aplikativni funktori - Applicatives

Uopštenje funktora na funkcije sa više argumenata

Definisati hijerarhiju za fmap, npr.

```
fmap0 :: a -> f a
fmap1 :: (a -> b) -> f a -> f b
fmap2 :: (a -> b -> c) -> f a -> f b -> f c
fmap3 :: (a -> b -> c -> d) -> f a -> f b -> f c -> f d
...
```

Ili **currying**

```
pure :: a -> f a
(<*>) :: f (a -> b) -> f a -> f b
```

pure – konvertuje tip a u tip f a

<*> - uopšten oblik primene funkcije gde su funkcija argument, vrednosti argumenta i vrednost rezultata su f strukture

Aplikativni funktori - Applicatives

```
pure g <*> x1 <*> x2 <*> ... <*> xn
```

Svaki argument x_i ima tip $f\ a_i$, a rezultat je tipa $f\ b$

```
fmap0 :: a -> f a
fmap0 = pure

fmap1 :: (a -> b) -> f a -> f b
fmap1 g x = pure g <*> x

fmap2 :: (a -> b -> c) -> f a -> f b -> f c
fmap2 g x y = pure g <*> x <*> y

fmap3 :: (a -> b -> c -> d) -> f a -> f b -> f c -> f d
fmap3 g x y z = pure g <*> x <*> y <*> z

...
```

```
class Functor f => Applicative f where
  pure :: a -> f a
  (<*>) :: f (a -> b) -> f a -> f b
```

Aplikativni funktori - Applicatives

Klasa funktora koja podržava `pure` i `<*>` naziva se aplikativni funktor

```
class Functor f => Applicative f where
  pure :: a -> f a
  (<*>) :: f (a -> b) -> f a -> f b
```

```
ghci> pure (+1) <*> Just 1
Just 2
ghci> pure (+) <*> Just 1 <*> Just 2
Just 3
ghci> pure (+) <*> Nothing <*> Just 2
Nothing
```

```
ghci> pure (\x -> \y -> \z -> x*y*z) <*> Just 1 <*> Just 2 <*> Just 3
Just 6
```

```
ghci> pure (*) <*> [1,2] <*> [3,4]
[3,4,6,8]
```

Aplikativni zakoni

```
pure id <*> x    = x
pure (g x)      = pure g <*> pure x
x <*> pure y     = pure (\g -> g y) <*> x
x <*> (y <*> z)  = (pure (.) <*> x <*> y) <*> z
```

Monadi

```
data Expr = Val Int | Div Expr Expr
eval :: Expr -> Int
eval (Val n) = n
eval (Div x y) = eval x `div` eval y
```

```
eval :: Expr -> Maybe Int
eval (Val n) = Just n
eval (Div x y) = case eval x of
  Nothing -> Nothing
  Just n -> case eval y of
    Nothing -> Nothing
    Just m -> safediv n m
```

```
ghci> eval (Div (Val 1) (Val 0))
*** Exception: divide by zero
```

```
safediv :: Int -> Int -> Maybe Int
safediv _ 0 = Nothing
safediv n m = Just (n `div` m)
```

```
ghci> eval (Div (Val 1) (Val 0))
Nothing
```

Monadi

Rešenje preko aplikativnih funktora?

```
eval :: Expr -> Maybe Int
eval (Val n) = pure n
eval (Div x y) = pure safediv <*> eval x <*> eval y
```

Neodgovarajući tipovi!!!

```
(>>==) :: m a -> (a -> m b) -> m b
mx >>= f = case mx of
    Nothing -> Nothing
    Just x -> f x
```

```
eval :: Expr -> Maybe Int
eval (Val n) = Just n
eval (Div x y) = eval x >>= \n ->
    eval y >>= \m ->
    safediv n m
```

Monadi

Uopšten oblik

```
m1 >>= \x1 ->  
m2 >>= \x2 ->  
...  
mn >>= \xn ->  
f x1 x2 ... xn
```

```
class Applicative m => Monad m where  
  return :: a -> m a  
  (>>=) :: m a -> (a -> m b) -> m b  
  
  return = pure
```

IO akcije, funkcije nad stablima, izrazima, do notacija, Maybe, ...

Zakoni za monade

```
return x >>= f      = f x
mx >>= return      = mx
(mx >>= f) >>= g    = mx >>= (\x -> (f x >>=g))
```

Tipovi i klase tipova

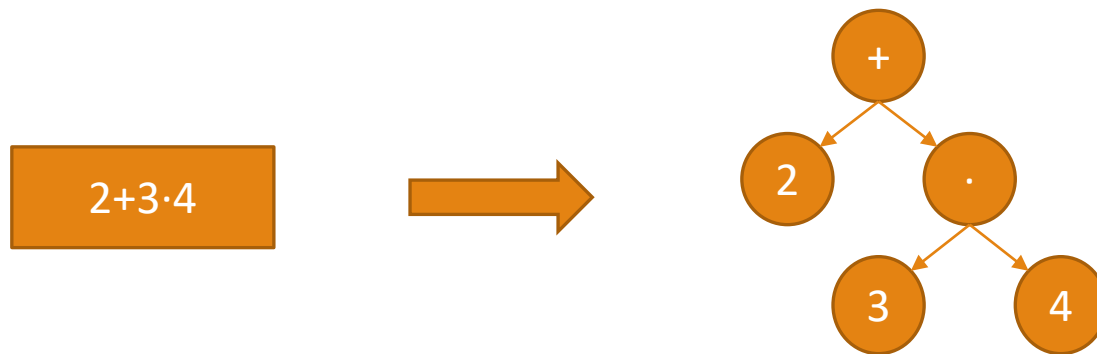
Type	Typeclasses
Bool	Eq, Ord, Show, Read, Enum, Bounded
Char	Eq, Ord, Show, Read, Enum, Bounded
Int	Eq, Ord, Show, Read, Enum, Bounded, Num, Real, Integral
Integer	Eq, Ord, Show, Read, Enum, Num, Real, Integral
Float	Eq, Ord, Show, Read, Enum, Num, Real, Fractional, RealFrac, Floating, RealFloat
Double	Eq, Ord, Show, Read, Enum, Num, Real, Fractional, RealFrac, Floating, RealFloat
Word	Eq, Ord, Show, Read, Enum, Bounded, Num, Real, Integral
Ordering	Eq, Ord, Show, Read, Enum, Bounded, Semigroup, Monoid
()	Eq, Ord, Show, Read, Enum, Bounded, Semigroup, Monoid
Maybe a	Eq, Ord, Show, Read, Semigroup, Monoid, Functor, Applicative, Monad, Foldable, Traversable
[a]	Eq, Ord, Show, Read, Semigroup, Monoid, Functor, Applicative, Monad, Foldable, Traversable
(a,b)	Eq, Ord, Show, Read, Bounded, Semigroup, Monoid, Functor, Applicative, Monad, Foldable, Traversable
a->b	Semigroup, Monoid, Functor, Applicative, Monad
IO	Semigroup, Monoid, Functor, Applicative, Monad
IOError	Eq, Show

Parser



Parser

Program koji analizira tekst i određuje njegovu sintaksnu strukturu



Svaki kompajler, interpreter, OS, veb čitač,... mora imati parser

```
type Parser = String -> Tree
```

Ovako definisan parser je “Black Box”

```
type Parser = String -> (Tree, String)
```

Vraća drvo i deo ulaza koji nije obrađen

Parser

Šta ako ulaz nije moguće parsirati u slučaju greške?

```
data Maybe          ~> [ ]  
  | Just a         ~> [a]
```

```
type Parser = String -> [(Tree, String)]
```

Različiti parseri će vratiti različite tipove stabla ili neku drugu vrednost

```
type Parser a = String -> [(a, String)]
```

Osnovni parser

```
import Control.Applicative
import Data.Char
```

Da bismo Parser koristili kao instancu klase korišćemo dummy konstruktor

```
newtype Parser a = P (String -> [(a,String)])
```

Ovakav parser se primenjuje na ulazni string korišćenjem funkcije koja uklanja konstruktor

```
parse :: Parser a -> String -> [(a,String)]
parse (P p) inp = p inp
```

Osnovni parser

Prvi karakter ulaza

```
item :: Parser Char
item = P (\inp -> case inp of
          [] -> []
          (x:xs) -> [(x,xs)])
```

```
ghci> parse item ""
[]
ghci> parse item "abc"
[('a',"bc")]
```

Parseri sekvenci

Parser kao instanca funktora

```
instance Functor Parser where
  -- fmap :: (a -> b) -> Parser a -> Parser b
  fmap g p = P (\inp -> case parse p inp of
                        [] -> []
                        [(v,out)] -> [(g v, out)])
```

```
ghci> parse (fmap toUpper item) "abc"
[('A',"bc")]
ghci> parse (fmap toUpper item) ""
[]
```

Parseri sekvenci

Parser kao instanca aplikativnog funktora

```
instance Applicative Parser where
  -- pure :: a -> Parser a
  pure v = P (\inp -> [(v,inp)])
  -- <*> :: Parser (a -> b) -> Parser a -> Parser b
  pg <*> px = P (\inp -> case parse pg inp of
    [] -> []
    [(g,out)] -> parse (fmap g px) out)
```

```
three :: Parser (Char,Char)
three = pure g <*> item <*> item <*> item
  where g x y z = (x,z)
```

```
ghci> parse (pure 1) "abc"
[(1,"abc")]
ghci> parse three "abcdef"
[(('a','c'),"def")]
ghci> parse three "ab"
[]
```

Parseri sekvenci

Parser kao instanca aplikativnog monada

```
instance Monad Parser where
  -- (>>=) :: Parser a -> (a -> Parser b) -> Parser b
  p >>= f = P (\inp -> case parse p inp of
    [] -> []
    [(v,out)] -> parse (f v) out)
```

```
three' :: Parser (Char,Char)
three' = do
  x <- item
  item
  z <- item
  return (x,z)
```

Donešenje odluka

Donešenje odluka nije specifično za parser, ali može da se generalizuje aplikativne tipove

```
class Applicative f => Alternative f where
  empty :: f a
  (<|>) :: f a -> f a -> f a
```

```
instance Alternative Maybe where
  empty :: Maybe a
  empty = Nothing
```

```
(<|>) :: Maybe a -> Maybe a -> Maybe a
Nothing <|> my = my
(Just x) <|> _ = Just x
```

Donešenje odluka

```
instance Alternative Parser where
  -- empty :: Parser a
  empty = P (\inp -> [])
  -- (<|>) :: Parser a -> Parser a -> Parser a
  p <|> q = P (\inp -> case parse p inp of
    [] -> parse q inp
    [(v,out)] -> [(v,out)])
```

```
ghci> parse empty "abc"
[]
ghci> parse (item <|> return 'd') "abc"
[('a',"bc")]
ghci> parse (empty <|> return 'd') "abc"
[('d',"abc")]
```

Izvedeni tipovi

```
sat :: (Char -> Bool) -> Parser Char
sat p = do
    x <- item
    if p x then return x else empty
```

```
digit :: Parser Char
digit = sat isDigit
```

```
lower :: Parser Char
lower = sat isLower
```

```
upper :: Parser Char
upper = sat isUpper
```

```
letter :: Parser Char
letter = sat isAlpha
```

```
alphanum :: Parser Char
alphanum = sat isAlphaNum
```

```
char :: Char -> Parser Char
char x = sat (== x)
```

Izvedeni tipovi

```
string :: String -> Parser String
string [] = return []
string (x:xs) = do
    char x
    string xs
    return (x:xs)
```

```
ghci> parse (char 'a') "abc"
[('a',"bc")]
ghci> parse (string "abc") "abcdef"
[("abc","def")]
ghci> parse (string "abc") "ab1234"
[]
```

```
ghci> parse (many digit) "123abc"
[("123","abc")]
ghci> parse (many digit) "abc"
[("", "abc")]
ghci> parse (some digit) "abc"
[]
```

Izvedeni tipovi

```
ident :: Parser String
ident = do
    x <- lower
    xs <- many alphanum
    return (x:xs)
```

```
nat :: Parser Int
nat = do
    xs <- some digit
    return (read xs)
```

```
ghci> parse ident "abc def"
[("abc", " def")]
ghci> parse nat "123 abc"
[(123, " abc")]
```

Izvedeni tipovi

```
space :: Parser ()
space = do
    many (sat isSpace)
    return ()
```

```
int :: Parser Int
int = do
    char '-'
    n <- nat
    return (-n)
<|> nat
```

```
ghci> parse space " abc"
[((), "abc")]
ghci> parse int "-123 abc"
[(-123, " abc")]
```

Obrada belina

```
token :: Parser a -> Parser a
token p = do
    space
    v <- p
    space
    return v
```

```
identifier :: Parser String
identifier = token ident
```

```
natural :: Parser Int
natural = token nat
```

```
integer :: Parser Int
integer = token int
```

```
symbol :: String -> Parser String
symbol xs = token (string xs)
```

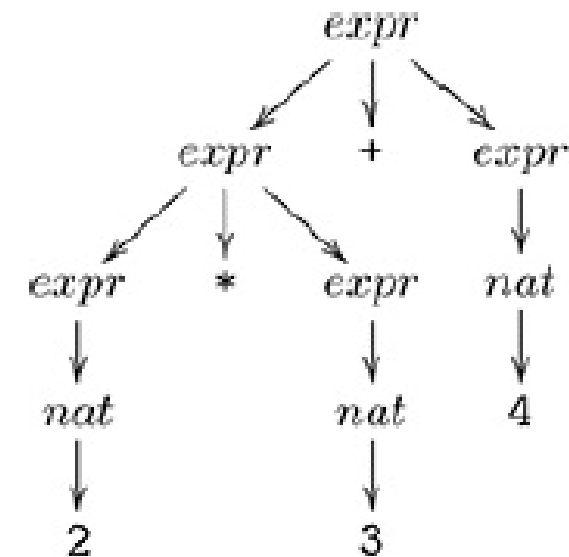
Aritmetički izrazi

$\text{expr} ::= \text{expr} + \text{expr} \mid \text{expr} * \text{expr} \mid (\text{expr}) \mid \text{nat}$

$\text{nat} ::= 0 \mid 1 \mid 2 \mid \dots$

Ovako definisana gramatika ne uzima u obzir činjenicu da je $*$ većeg prioriteta od $+$

$2 * 3 + 4$



Aritmetički izrazi

$\text{expr} ::= \text{expr} + \text{expr} \mid \text{term}$

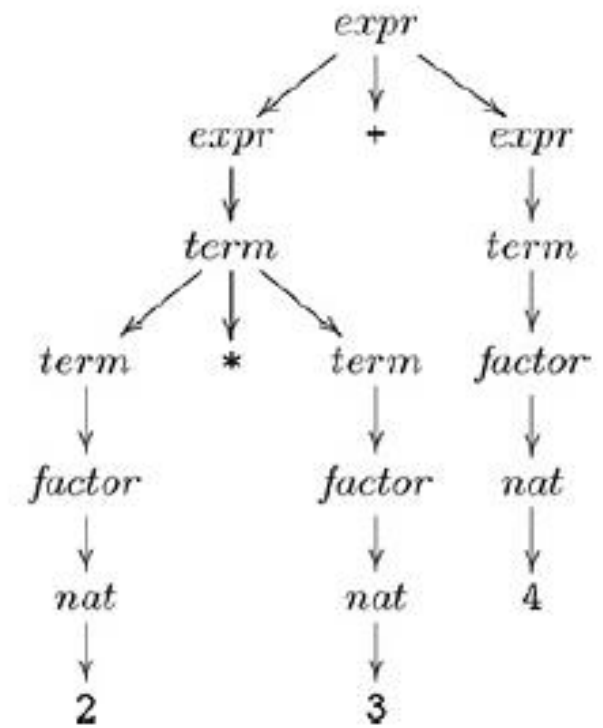
$\text{term} ::= \text{term} * \text{term} \mid \text{factor}$

$\text{factor} ::= (\text{expr}) \mid \text{nat}$

$\text{nat} ::= 0 \mid 1 \mid 2 \mid \dots$

Prioritet je rešen, ali nije uzeta u obzir
desna asocijativnost za $*$ i $+$

$2 * 3 + 4$



Aritmetički izrazi

$\text{expr} ::= \text{term} + \text{expr} \mid \text{term}$

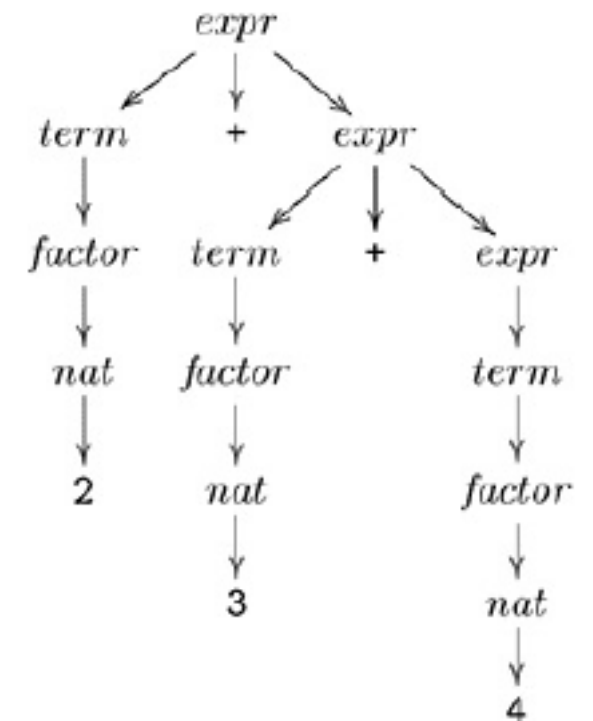
$\text{term} ::= \text{factor} * \text{term} \mid \text{factor}$

$\text{factor} ::= (\text{expr}) \mid \text{nat}$

$\text{nat} ::= 0 \mid 1 \mid 2 \mid \dots$

Dobijena gramatika je **nedvosmislena**

$2 + 3 + 4$



Aritmetički izrazi

`expr ::= term (+ expr | ϵ)`

`term ::= factor (* term | ϵ)`

`factor ::= (expr) | nat`

`nat ::= 0 | 1 | 2 | ...`

```
expr :: Parser Int
```

```
expr = do
```

```
    t <- term
```

```
    do
```

```
        symbol "+"
```

```
        e <- expr
```

```
        return (t + e)
```

```
    <|> return t
```

```
term :: Parser Int
```

```
term = do
```

```
    f <- factor
```

```
    do
```

```
        symbol "*"
```

```
        t <- term
```

```
        return (f * t)
```

```
    <|> return f
```

```
factor :: Parser Int
```

```
factor = do
```

```
    symbol "("
```

```
    e <- expr
```

```
    symbol ")"
```

```
    return e
```

```
<|> natural
```

Aritmetički izrazi

```
eval' :: String -> Int
eval' xs = case (parse expr xs) of
    [(n,[])] -> n
    [(_,out)] -> error ("Unused input " ++ out)
    [] -> error "Invalid input"
```

```
ghci> eval' "2*3+4"
10
ghci> eval' "2*(3+4)"
14
```

```
ghci> eval' "2*3^4"
*** Exception: Unused input ^4
CallStack (from HasCallStack):
  error, called at parser.hs:176:27 in main:Main
ghci> eval' "one plus two"
*** Exception: Invalid input
CallStack (from HasCallStack):
  error, called at parser.hs:177:27 in main:Main
ghci>
```

Kalkulator

```
box :: [String]
box = ["+-----+",
      "|           |",
      "+---+---+---+---+",
      "| q | c | d | = |",
      "+---+---+---+---+",
      "| 1 | 2 | 3 | + |",
      "+---+---+---+---+",
      "| 4 | 5 | 6 | - |",
      "+---+---+---+---+",
      "| 7 | 8 | 9 | * |",
      "+---+---+---+---+",
      "| 0 | ( | ) | / |",
      "+---+---+---+---+"]
```

Kalkulator

```
buttons :: String
buttons = standard ++ extra
  where
    standard = "qcd=123+456-789*0()/"
    extra = "QCD \ESC\BS\DEL\n"
```

```
showbox :: IO ()
showbox = sequence_ [writeAt (1,y) b | (y,b) <- zip [1..] box]
```

```
display xs = do
  writeAt (3,2) (replicate 13 ' ')
  writeAt (3,2) (reverse (take 13 (reverse xs)))
```

Kalkulator

```
calc :: String -> IO ()
calc xs = do
    display xs
    c <- getCh
    if elem c buttons then
        process c xs
    else
        do
            beep
            calc xs
```

```
process :: Char -> String -> IO ()
process c xs | elem c "qQ\ESC" = quit
             | elem c "dD\BS\DEL" = delete xs
             | elem c "=\n" = eval xs
             | elem c "cC" = clear
             | otherwise = press c xs
```


Kalkulator

```
clear :: IO ()  
clear = calc []
```

```
press :: Char -> String -> IO ()  
press c xs = calc (xs ++ [c])
```

```
run :: IO ()  
run = do  
    cls  
    showbox  
    clear
```

